

CSE4213 Lecture Notes

Predicate Logic and Substitutions

Schneider, chapter 2

Computer Science and Software Engineering
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20050326 / Lecture 9

Outline

Predicate Calculus

Substitutions

Summary

Predicates

- **Predicates** are statements which are either true or false
- Predicates can be used to make statements about sets, and hence are very useful in B
- $x \in PRICE$ is a predicate that states something about x , a potential member of the set $PRICE$. If true, then x is an element of $PRICE$
- Hence use predicates to define sets with **set comprehension**
- $\{x \mid x \in \mathbb{N} \wedge x \leq 10\}$

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Operations on Predicates

- **negation** (not): $\neg P$
- **disjunction** (or): $P \vee Q$
- **conjunction** (and): $P \wedge Q$
- **implication** (implies): $P \Rightarrow Q$
- **equivalence** (if and only if): $P \Leftrightarrow Q$

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Quantification

- We often want to state predicates across a set
- **universal quantification** defines a predicate across all members of a set: $\forall x.(x \in S \Rightarrow P)$
- **existential quantification** defines a predicate for at least one member of a set: $\exists x.(x \in S \wedge P)$
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Duality of Quantification

- If universal quantification is false, then there must be at least one element that makes it false
- If existential quantification is false, then all elements must not satisfy the predicate.
- Hence we have:

$$\neg \forall x.(x \in S \Rightarrow P) \Leftrightarrow \exists x.(x \in S \wedge \neg P)$$

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Constraining Predicates

- There are some contexts where it is stated that a predicate P must **constrain** some list of variables z
- To constrain the variable x , the predicate P must contain predicates of the form: $x \in S$, $x \subseteq S$, $x \subset S$, or $x = E$, where $x \setminus S$, $x \setminus E$ and S is a set, and E is an expression.
- $x \setminus E$ means that x is not free in E
- We say that x is a **bound variable** in P .

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Free Variables

- A variable x is free in an expression if it is not bound by a quantifier.
- All free instances of z in P and Q are bound in $\forall z.(P \Rightarrow Q)$ and $\exists z.(P \wedge Q)$
- Hence, $z \setminus \forall z.(P \Rightarrow Q)$
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Substitutions

Substitutions as Assignments

- **Substitutions** are crucial to the B-Method
- Substitutions are the way in which we effect change of state
- A substitution (declarative) is the formal equivalent of assignment (imperative)
- $P[E/x]$ means: *P , with every occurrence of x replaced by E*
- multiple substitutions are possible: $P[E, F, \dots / x, y, \dots]$

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General Substitution Language

- The substitutions shown in the previous slide are part of the **General Substitution Language** (GSL), a formalism which we will not use much in this unit.
- Instead use the **Abstract Machine Notation**
- Think of substitutions as **Predicate Transformers**

$$[G]P$$

- Example: expanding $[x := 2]x < 5$ gives $2 < 5$ (a true predicate)
- Substitutions may or may not **satisfy** a predicate
- *The set of variables z satisfying the predicate P* means the variables z are instantiated to values that when substituted for free instances of z in P make the predicate true.

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- The philosophy of program specification may be stated thus:
- Have some starting situation (state) identified by predicate Q
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- Predicates define **states of interest**
- **Substitutions** define how values (state) change
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