Extension of universal cycles for globally identifying colorings of cycles

Pierre Coupechoux

LAAS-CNRS, France

Discrete Maths Research Group, November 21, 2016

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Identifying codes

Definition

An identifying code of a graph is a subset C of the vertices such that:

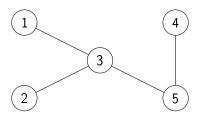
- Every vertex has a neighbor in C (Domination);
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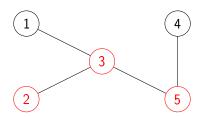


Identifying codes

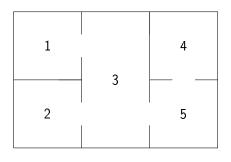
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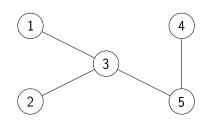
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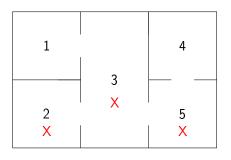


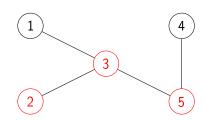
A classical application of an identifying code



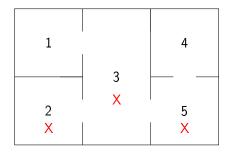


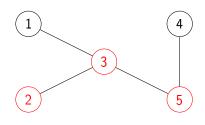
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A classical application of an identifying code





- Domination: each room is guarded.
- Identification: in case of fire, we know exactly where the fire is.

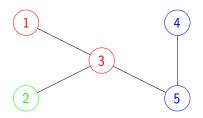
Variant: identifying coloring

- Each vertex has a color.
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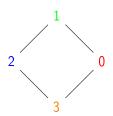
	Red	Green	Blue
1	Х		
2	Х	Х	
3	Х	Х	Χ
4			Χ
5	Х		Χ

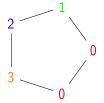
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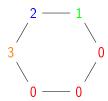
- Fixed number of colors L.
- Longest cycle with a globally identifying coloring.

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• The proof is constructive.



• A vertex: Up to 3 colors in its neighborhood.



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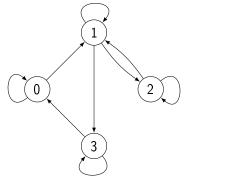
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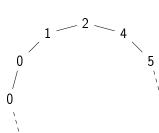


- A set of at most 3 colors: initial color + two differences.
- We arbitrary chose the order of the 3 colors.
 - For example, $\{0,1\}$ will only appear with colors 0, 0, and 1, in this order.

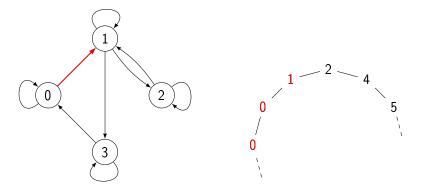
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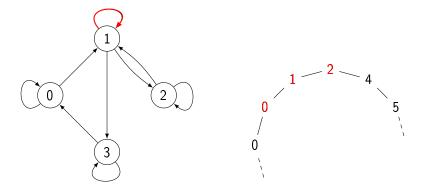




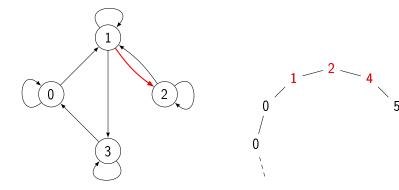
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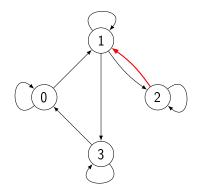
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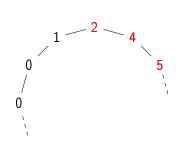


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• Edges (0, i) and $(i, 0) \longrightarrow$ same color sets.

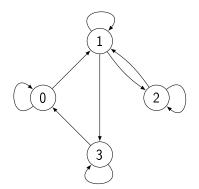
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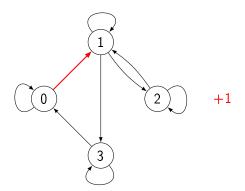
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- The graph has to be Eulerian.

- Choice of a circuit in the graph.
- Sequence of differences.

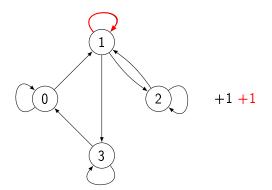
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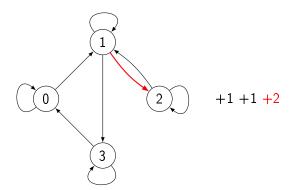
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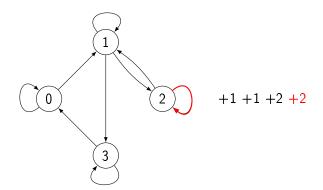
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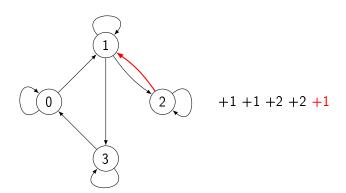
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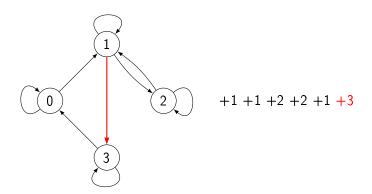
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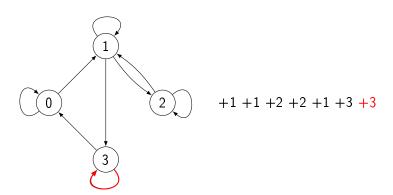
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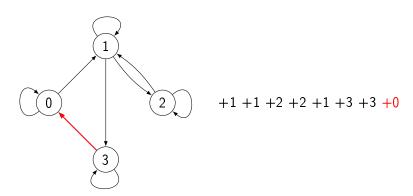
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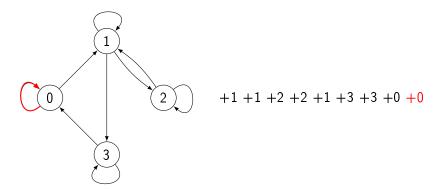
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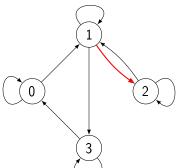
$$1 + 1 + 1$$
 $0 - 1 - 2$



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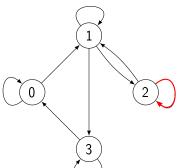
$$1 + 1 + 1 + 2$$

 $0 - 1 - 2 - 4$

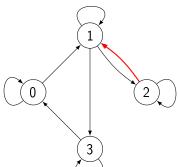


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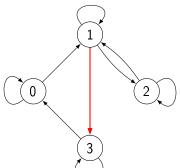
$$0 - 1 - 2 - 4 - 6$$



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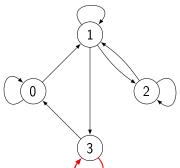


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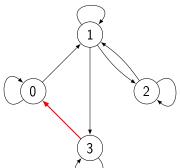


November 21, 2016

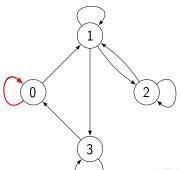
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+1 +1 +2 +2 +1 +3 +3 +0 +0
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0		1		2		4		6		7		2		5		5		5
5		6		7		1		3		4		7		2		2		2

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	+1	+1	+2	<u>)</u>	+2		+1		+3	+	-3	+0	+	0
0	1		2	4		6		7		2	5		5	5
5	6		7	1		3		4		7	2		2	2
2	3		4	6		0		1		4	7		7	7

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0	1	2	4	6	7	2	5	5	5
5	6	7	1	3	4	7	2	2	2
2	3	4	6	0	1	4	7	7	7
7	0	1	3	5	6	1	4	4	4

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5	6	7	1	3	4	7	2	2	2
2	3	4	6	0	1	4	7	7	7
7	0	1	3	5	6	1	4	4	4
4	5	6	0	2	3	6	1	1	1

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5	6	7	1	3	4	7	2	2	2
2	3	4	6	0	1	4	7	7	7
7	0	1	3	5	6	1	4	4	4
4	5	6	0	2	3	6	1	1	1
1	2	3	5	7	0	3	6	6	6

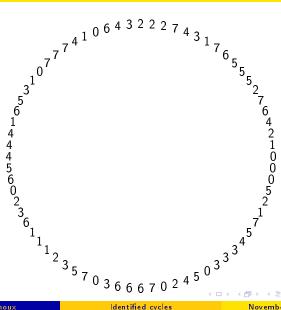
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0	1	2	4	6	7	2	5	5	5
5	6	7	1	3	4	7	2	2	2
2	3	4	6	0	1	4	7	7	7
7	0	1	3	5	6	1	4	4	4
4	5	6	0	2	3	6	1	1	1
1	2	3	5	7	0	3	6	6	6
6	7	0	2	4	5	0	3	3	3

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0	1		2	4		6		7		2		5		5		5
5	6		7	1		3		4		7		2		2		2
2	3		4	6		0		1		4		7		7		7
7	0		1	3		5		6		1		4		4		4
4	5		6	0		2		3		6		1		1		1
1	2		3	5		7		0		3		6		6		6
6	7		0	2		4		5		0		3		3		3
3	4		5	7		1		2		5		0		0		0
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Cycle with 8 colors



• What if we obtain 0 too soon ?

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• Lose vertex with initial color L-1 and differences +1, +1.

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- Add them at the end!
 - No missing (+1, +1) differences
 - Path starting and ending with $0 \approx \text{cycle}$

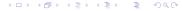


Improvement

- When L odd, we can add some vertices with differences $(0, \frac{L-1}{2}), (\frac{L-1}{2}, \frac{L-1}{2})$ and $(\frac{L-1}{2}, 0)$.
- Path starting and ending with two 0s.

$$0 - 0 - 4 - 8 - 3 - 7 - 2 - 6 - 1 - 5 - 0 - 0$$

Easy to insert in the previous cycle.



Results

With L colors, cycles build with this method have length:

$$\frac{L^{3} + 5L}{6} - \begin{cases} \frac{L+2}{4}L & \text{if } L \equiv 0[4] \\ \frac{L+3}{4}L - 1 & \text{if } L \equiv 1[4] \\ \frac{L^{2}}{4} & \text{if } L \equiv 2[4] \\ \frac{L+1}{4}L - 1 & \text{if } L \equiv 3[4] \end{cases} - \begin{cases} \frac{L}{3} & \text{if } L \equiv 0[3] \\ 0 & \text{else} \end{cases}$$



Not optimal

- For L=5, we can build a cycle with $\frac{L^3+5L}{6}=25$ vertices.
 - Choosing differences (+1,+2) instead of (+2,+2).
- Not for L = 4 (optimal = 10 < 14).



Thank you (in French)

