

# A survey of Tutte-Whitney polynomials

Graham Farr

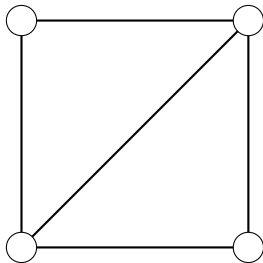
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July 2007

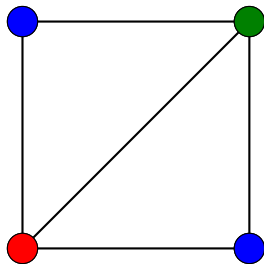
# Counting colourings

- ▶ proper colourings



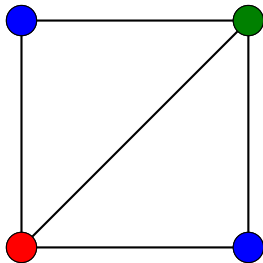
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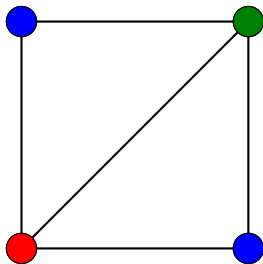
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Adjacent vertices receive  
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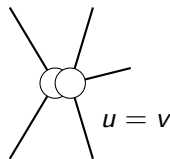
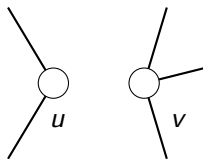
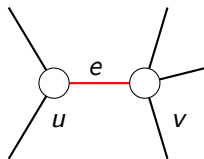
- ▶ chromatic polynomial:

$$P(G; q) = \# \text{ } q\text{-colourings of } G$$

# Deletion-contraction

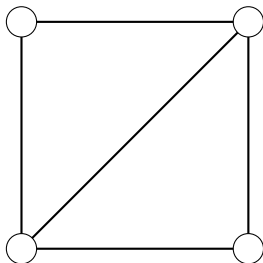
For any edge  $e$ :

$$P(G; q) = P(G \setminus e; q) - P(G/e; q)$$



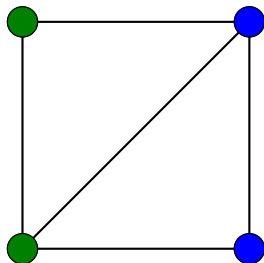
## Partition functions: Potts models

- ▶ general  $q$ -colourings (may be improper)



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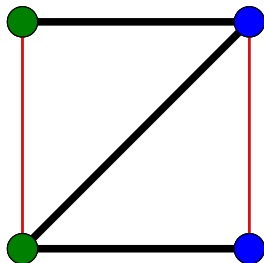
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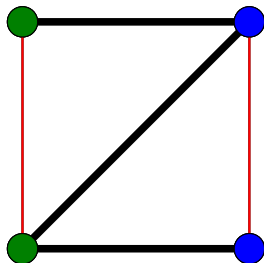
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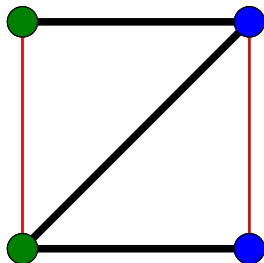
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**Good** and **bad** edges

# Partition functions: Potts models

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**Good** and **bad** edges

- ▶ Partition function:

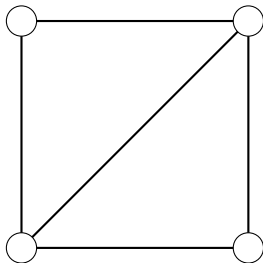
$$Z(G; K, q) = \sum_{\substack{\text{all } q\text{-colourings} \\ \text{(not just proper)}}} e^{-K \cdot (\# \text{ good edges})}$$

## All-terminal reliability

- ▶ Choose edges randomly:  $\Pr(\text{edge}) = p$

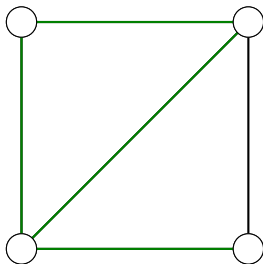
## All-terminal reliability

- ▶ Choose edges randomly:  $\Pr(\text{edge}) = p$
- ▶ Want chosen edges to contain a *spanning tree*



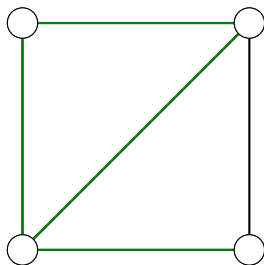
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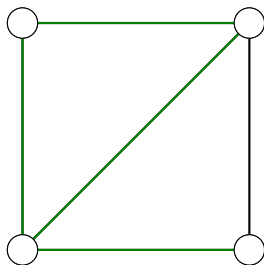
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chosen edges

- ▶ Reliability:

$$\Pi(G, p) = \Pr(\text{chosen edges contain a spanning tree})$$



... etc

... etc

- ▶ flow polynomial

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- ▶ # spanning trees, forests, spanning subgraphs

... etc

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# Tutte-Whitney polynomials

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for all  $X \subseteq E$ :

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- ▶ Tutte polynomial:

$$T(G; x, y) = R(G; x - 1, y - 1).$$

# The “Recipe Theorem”

## Theorem

(Tutte 1947  $\rightarrow$  Brylawski 1972  $\rightarrow$  Oxley & Welsh 1979)

If a function  $f$  on graphs ...

- ▶ is *invariant* under isomorphism,
- ▶ satisfies a *deletion-contraction* relation,
- ▶ is *multiplicative* over components  
(i.e.,  $f(G_1 \cup G_2) = f(G_1) \cdot f(G_2)$ ),

... then  $f$  is essentially a (partial) evaluation of the Tutte-Whitney polynomial.

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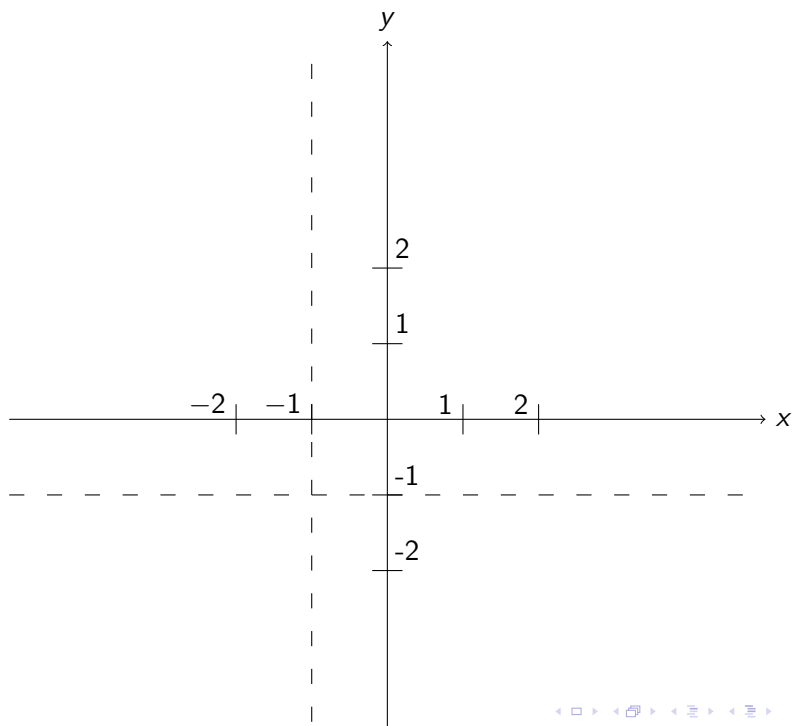
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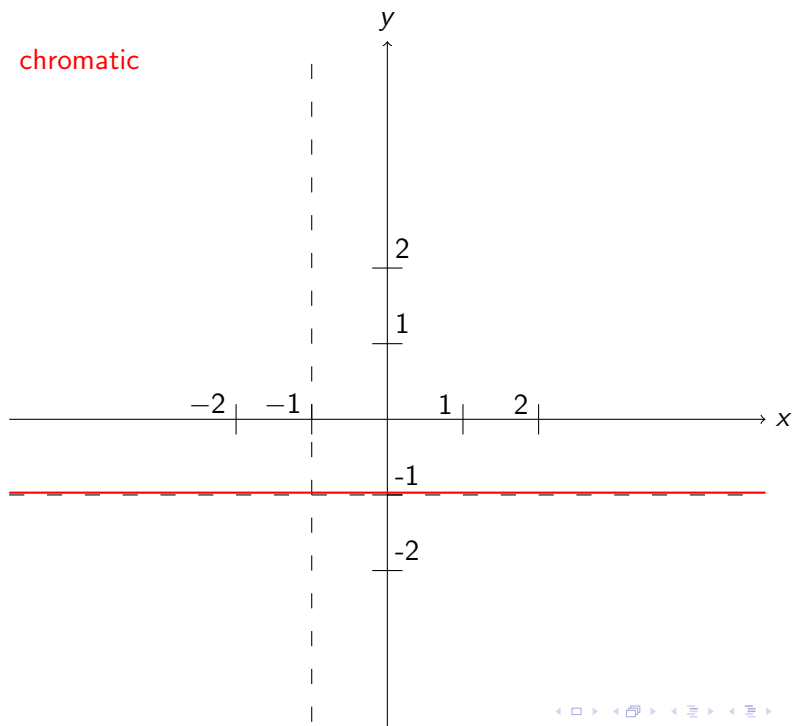
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## Example

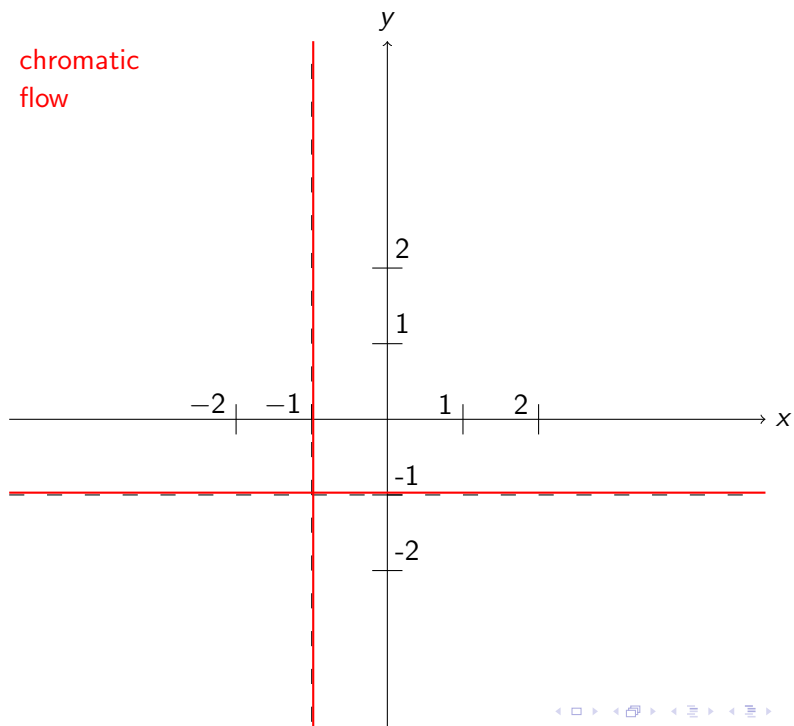
$$P(G; q) = (-1)^{\rho(E)} q^{k(G)} R(G; -q, -1)$$



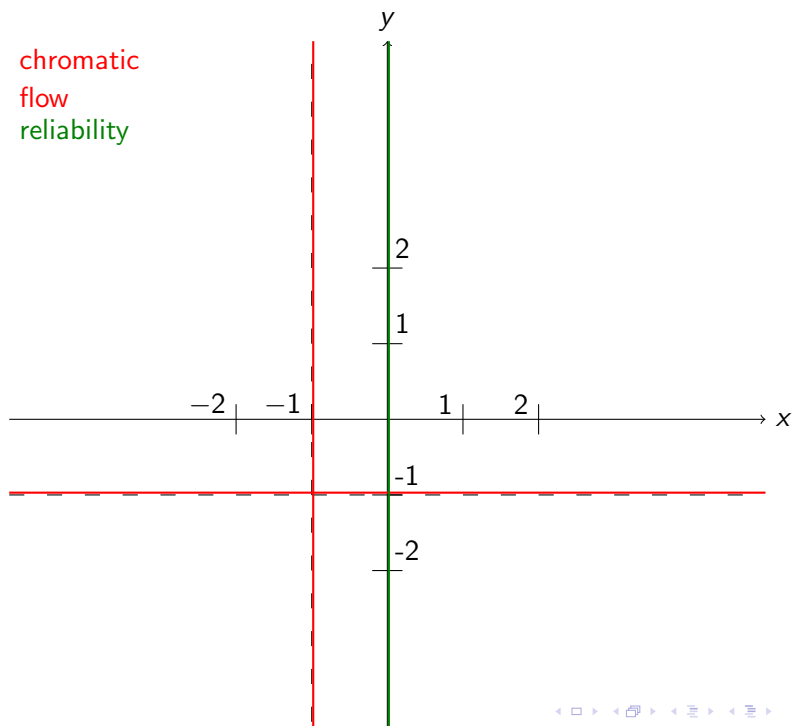
chromatic



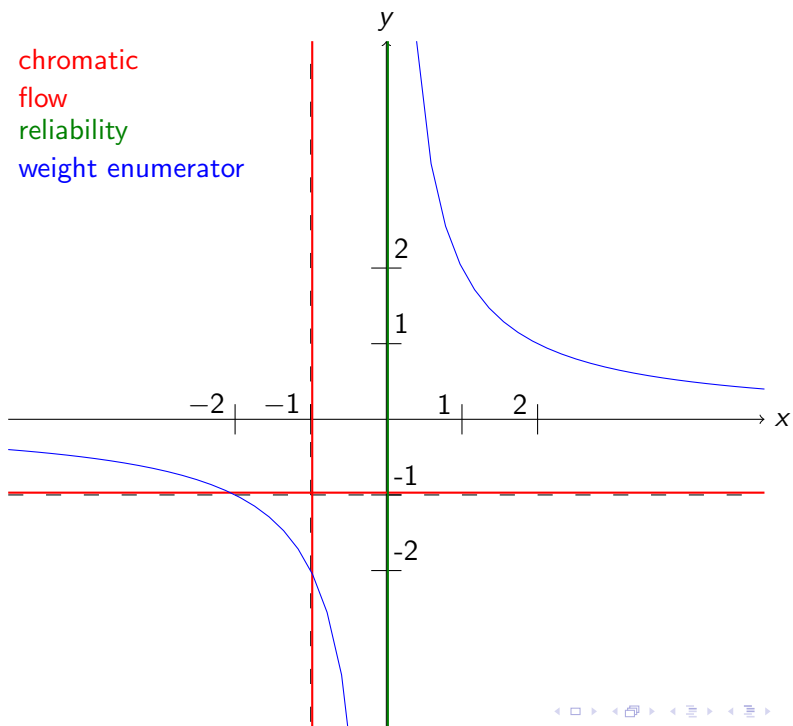
chromatic  
flow



chromatic  
flow  
reliability

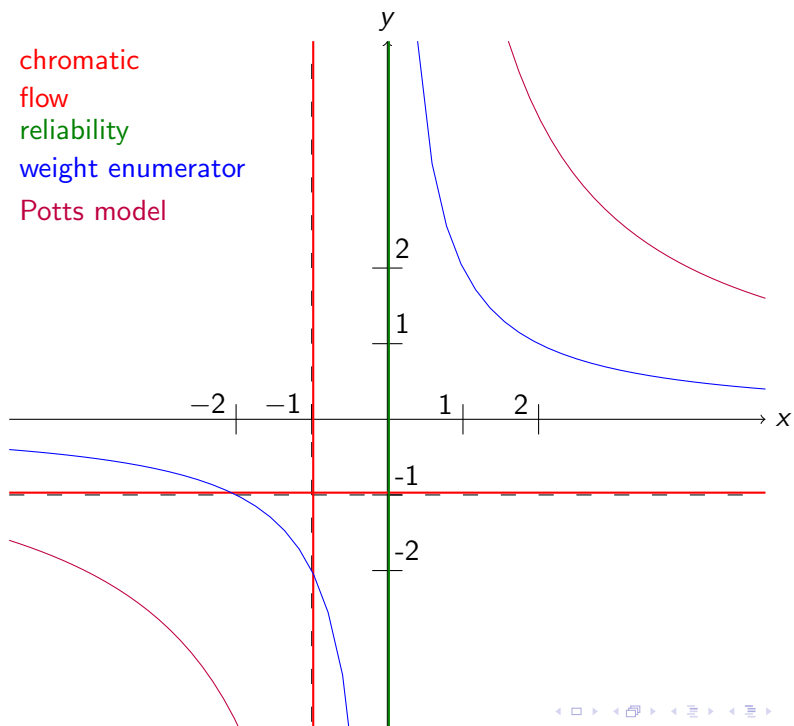


chromatic  
flow  
reliability  
weight enumerator

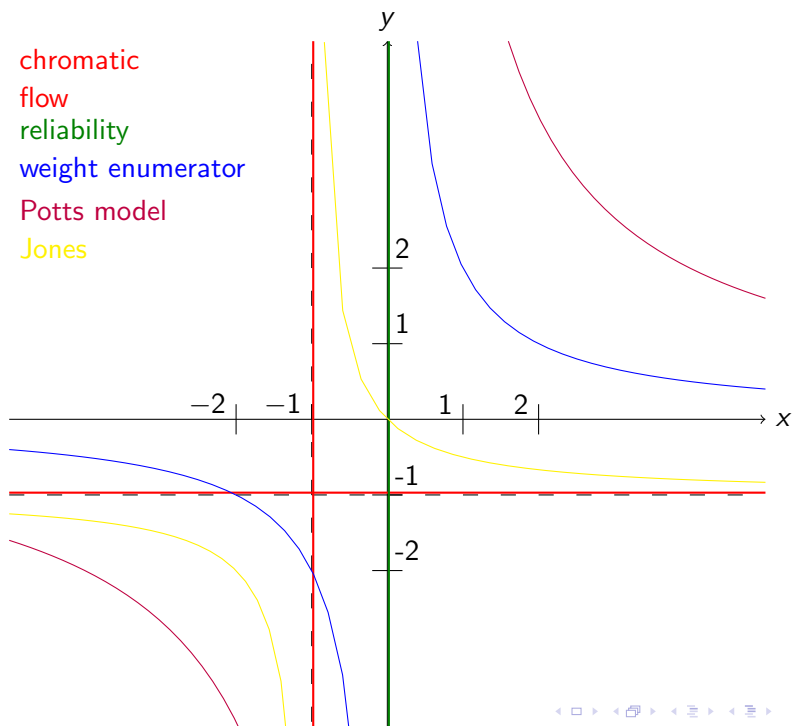




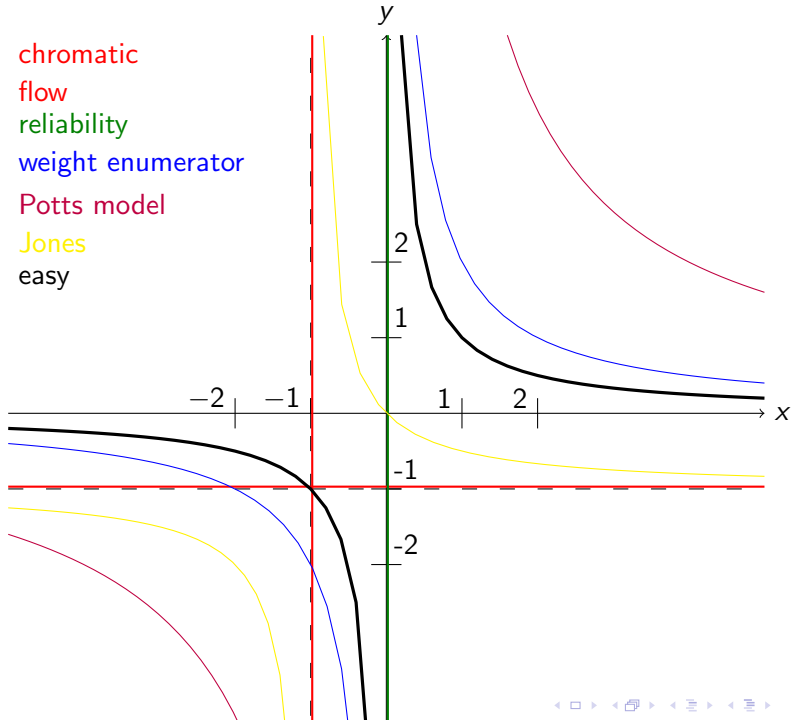
chromatic  
flow  
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Potts model



chromatic  
flow  
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Jones



chromatic  
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Jones  
easy



chromatic

flow

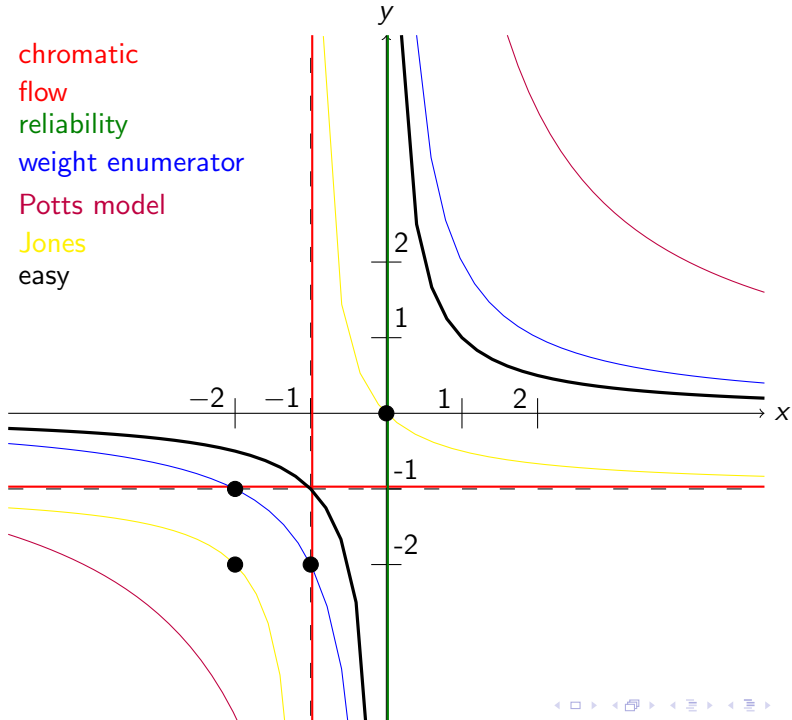
reliability

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# History

# History

Graphs:  
Chrom.  
poly

# History

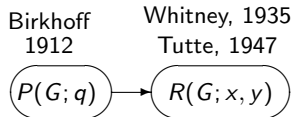
Graphs:  
Chrom.  
poly

Birkhoff  
1912

$$P(G; q)$$

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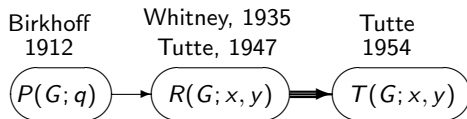
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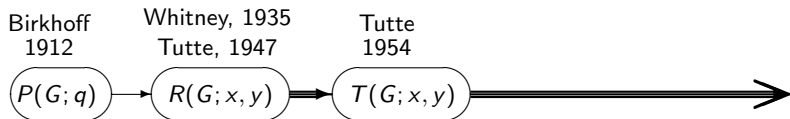
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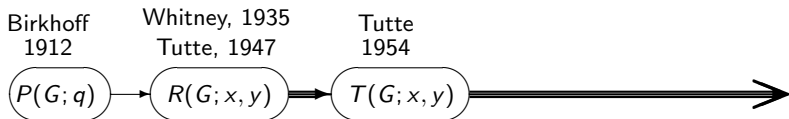
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# History

Stat Mech:  
partition  
functions

Graphs:  
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Ising  
1925

Ising  
model  
( $q = 2$ )

Graphs:  
Chrom.  
poly

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1912

$P(G; q)$

Whitney, 1935  
Tutte, 1947

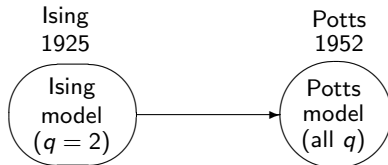
$R(G; x, y)$

Tutte  
1954

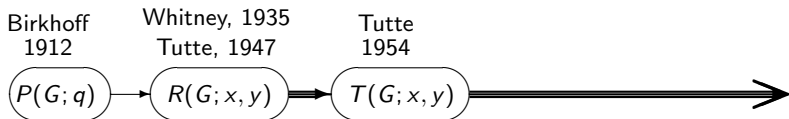
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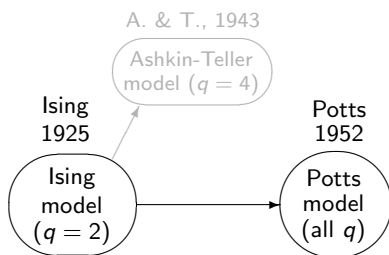


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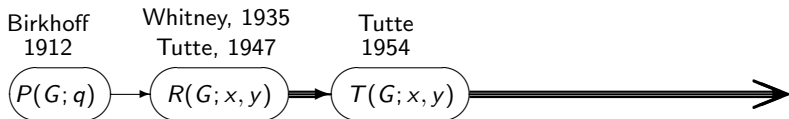


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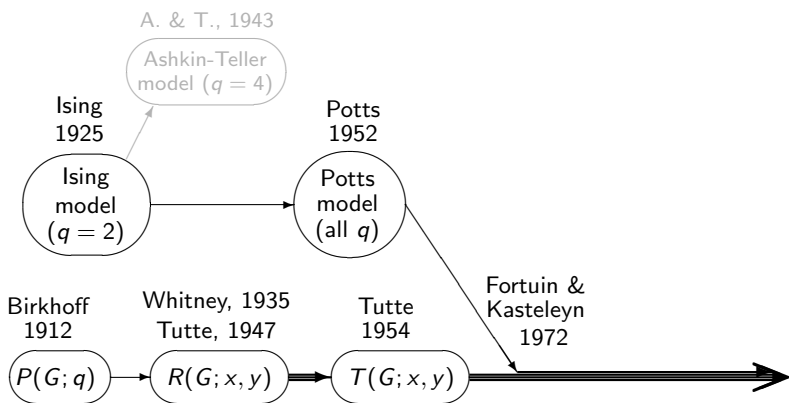
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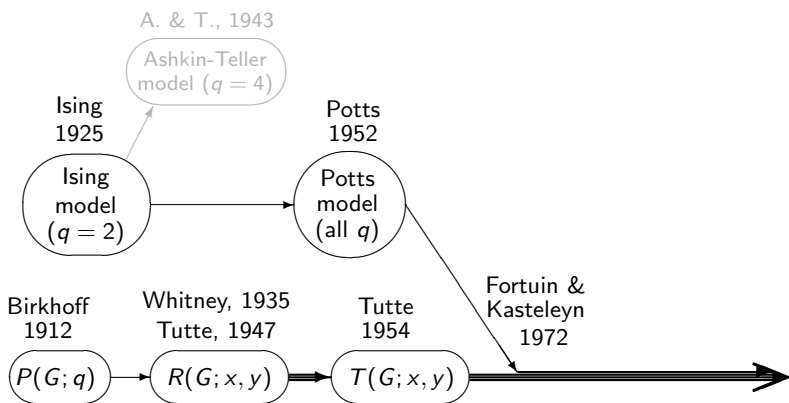


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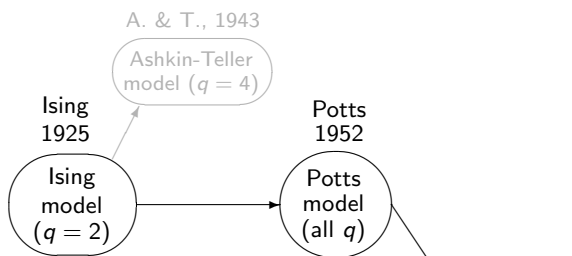
Linear codes:  
weight  
enumerator



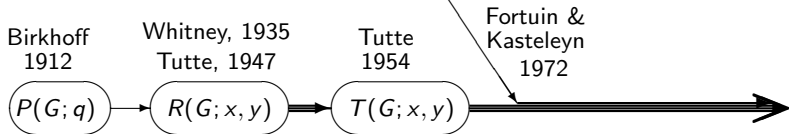


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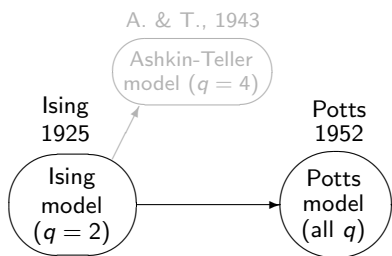
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MacWilliams  
1963

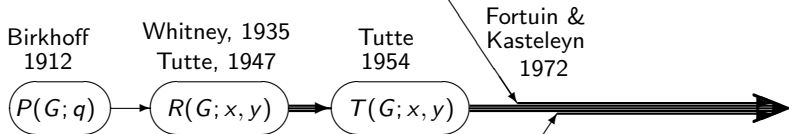
$$A_C(z)$$

# History

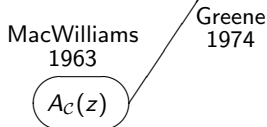
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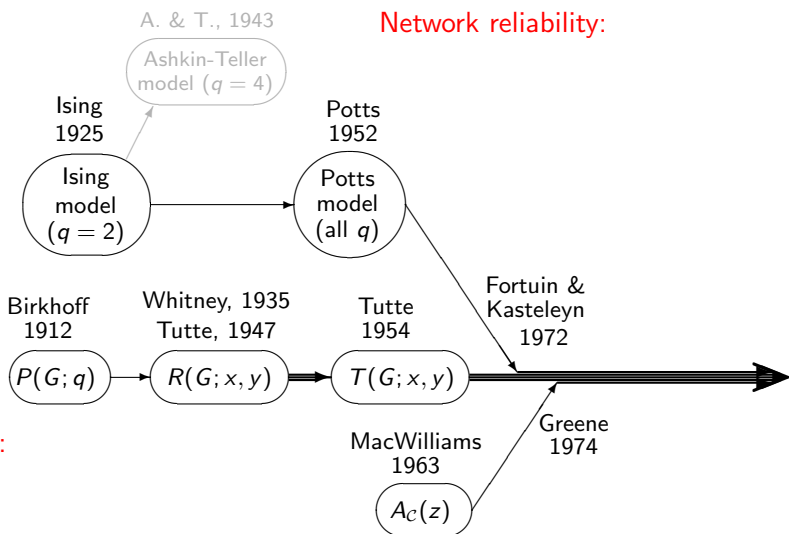
# History

## Network reliability:

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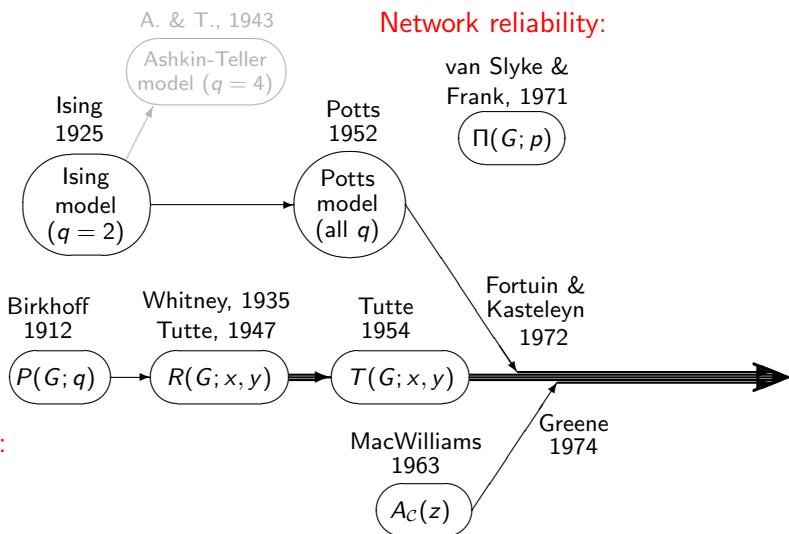
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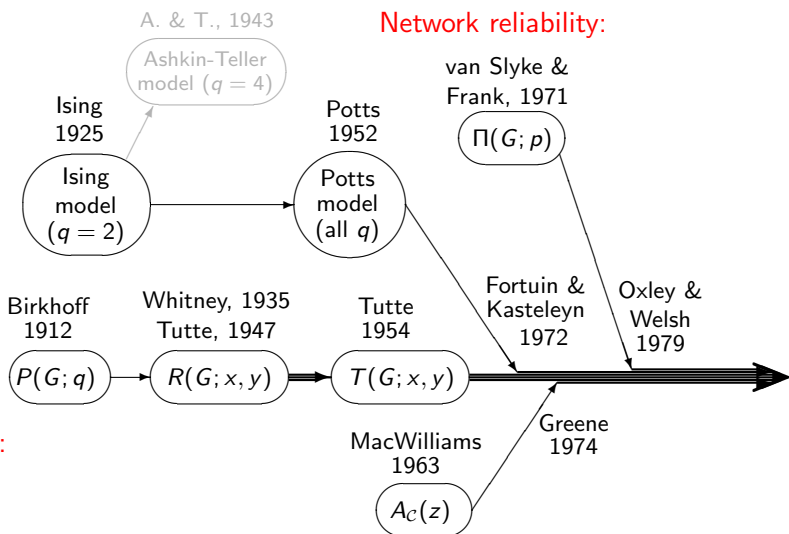
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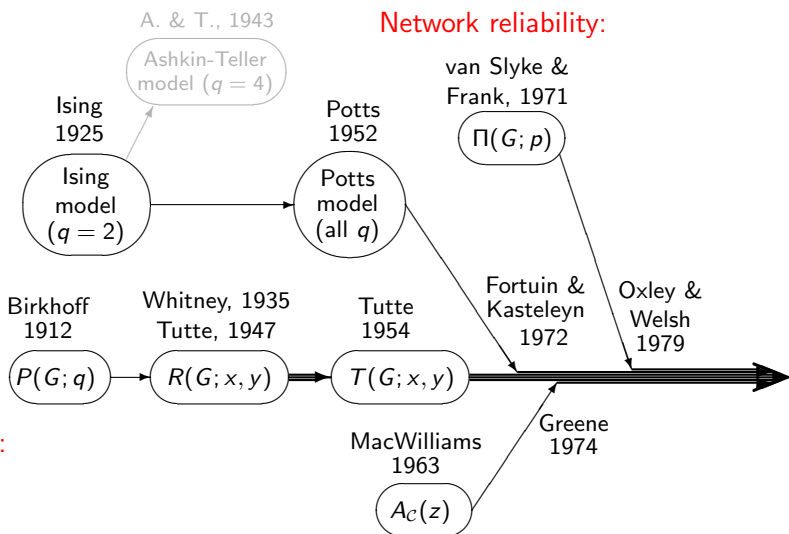
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Knots:  
Jones poly

Network reliability:



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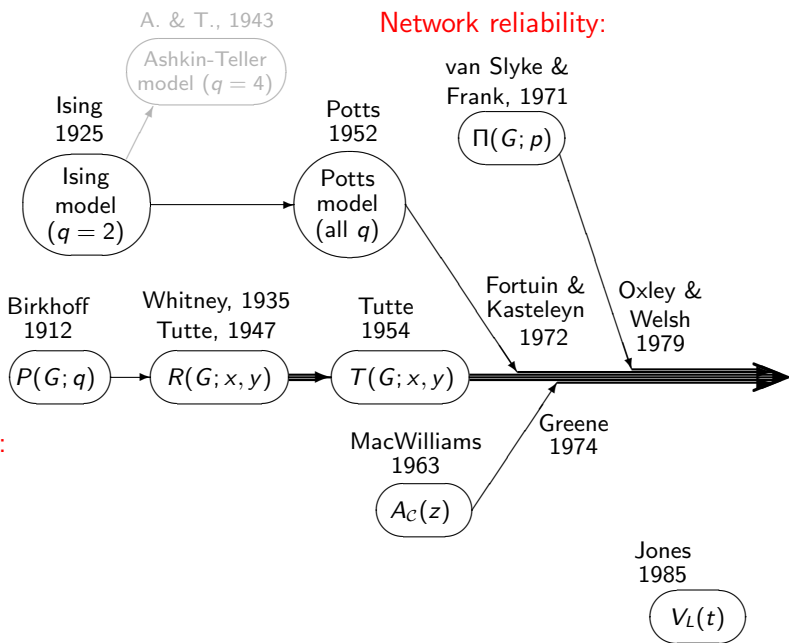
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# Complexity of computing all of $R(G; x, y)$

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- ▶ Square grid subgraphs, max deg 3:  $\#P$ -hard (GF, 2006)
- ▶ Square grid graphs: Open (in  $\#P_1$ )
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# Complexity of evaluating at specific points

## Theorem

*(Jaeger, Vertigan and Welsh, 1990)*

*The problem of determining  $R(G; x, y)$ , given a graph  $G$ , is  $\#P$ -hard at all points  $(x, y)$  except those where  $xy = 1$  and  $(x, y) = (0, 0), (-1, -2), (-2, -1), (-2, -2)$ .*



# Generalisations

Extensions from **graphs** to:

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- ▶ representable matroids (Smith), matroids (Tutte, Crapo), greedoids (Gordon & McMahon), Boolean functions or set systems (GF), hyperplane arrangements (Welsh & Whittle, Ardila), semimatroids (Ardila), signed graphs (Murasugi), rooted graphs (Wu, King & Lu),  $K$ -terminal graphs (Traldi), biased graphs (Zaslavsky), matroid perspectives (Las Vergnas), matroid pairs (Welsh & Kayibi), bimatroids (Kung), graphic polymatroids (Borzacchini), general polymatroids (Oxley & Whittle), ...

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... or extend the **polynomials**:

- ▶ multivariate polynomials of various kinds: variables at each **vertex** (Noble & Welsh), or **edge** (Fortuin & Kasteleyn, Traldi, Kung, Sokal, Bollobás & Riordan, Zaslavsky, Ellis-Monaghan & Riordan, Britz).

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- ▶ interesting partial evaluations
- ▶ deletion-contraction relations
- ▶ Recipe Theorems
- ▶ easier proofs
- ▶ roots
- ▶ how much of the graph is determined by the polynomial?

We now look at a generalisation to Boolean functions . . .

Rank  $\leftrightarrow$  rowspace

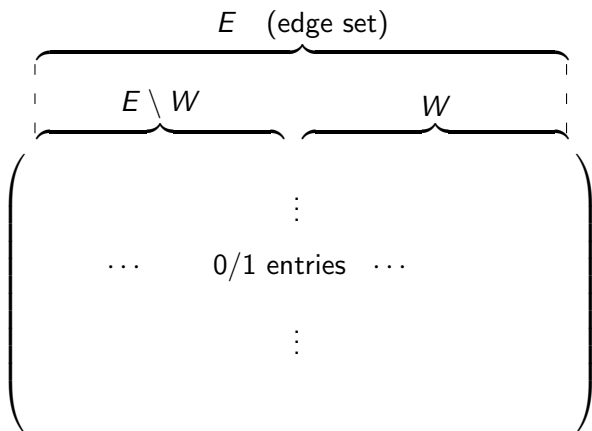
Incidence matrix

$$\begin{array}{c} \text{edges} \\ \text{vertices} \end{array} \left( \begin{array}{ccc} & \vdots & \\ \cdots & 0/1 \text{ entries} & \cdots \\ & \vdots & \end{array} \right)$$

Rank  $\leftrightarrow$  rowspace

Incidence matrix

vertices

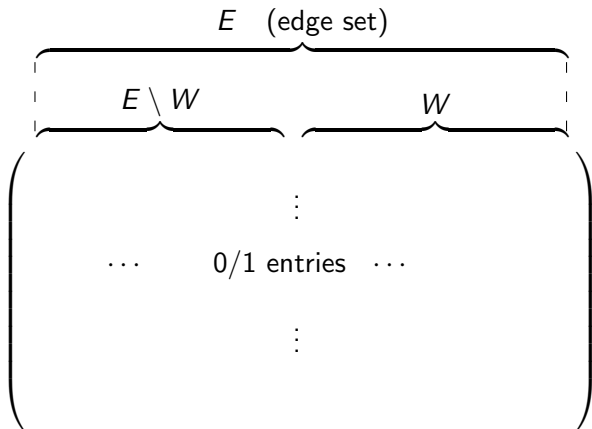


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$\longrightarrow$  echelon form

vertices





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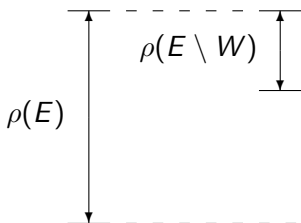
$\longrightarrow$  echelon form

$$\begin{array}{c} \overbrace{\hspace{10em}}^{E \text{ (edge set)}} \\ \overbrace{\hspace{4em}}^{E \setminus W} \quad \overbrace{\hspace{4em}}^W \\ \left( \begin{array}{c|c|c|c|c} 0 & I & \dots & \dots & \dots \\ \hline & & & & \\ \hline & 0 & & 0 & I & \vdots \\ \hline & & & 0 & & \end{array} \right) \end{array}$$

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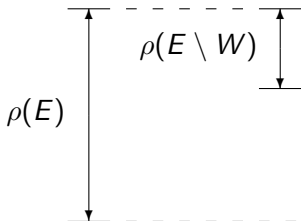


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$$(Q^\dagger \rho)(V) = (-1)^{|V|} \sum_{W\subseteq V} (-1)^{|W|} 2^{\rho(E)-\rho(E\setminus W)}$$

# Properties of the transform $Q$

Basic properties:

- ▶  $(Q^\dagger Q f)(V) = \frac{f(V)}{f(\emptyset)}$
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	$f$	$\xrightarrow{Q}$	$Qf$
Hadamard transform	$\downarrow$		$\downarrow$ matroid-style dual
	$\hat{f}$	$\xrightarrow{Q}$	$(Qf)^* = Q\hat{f}$

## Extending the Whitney rank generating function

$$R(f; x, y) = \sum_{X \subseteq E} x^{Qf(E) - Qf(X)} y^{|X| - Qf(X)}$$

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Example:

$$E = \{1, 2\}$$

$X$	$f$
$\emptyset$	1
$\{1\}$	1
$\{2\}$	1
$\{1, 2\}$	0



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$$R(f; x, y) = x^{\log_2 3} + 2xy^{2 - \log_2 3} + y^{2 - \log_2 3}$$

# Deletion-contraction

For  $e \in E$ ,  $X \subseteq E \setminus \{e\}$ :

Deletion

$$(f \parallel e)(X) = \frac{f(X) + f(X \cup \{e\})}{f(\emptyset) + f(\{e\})};$$

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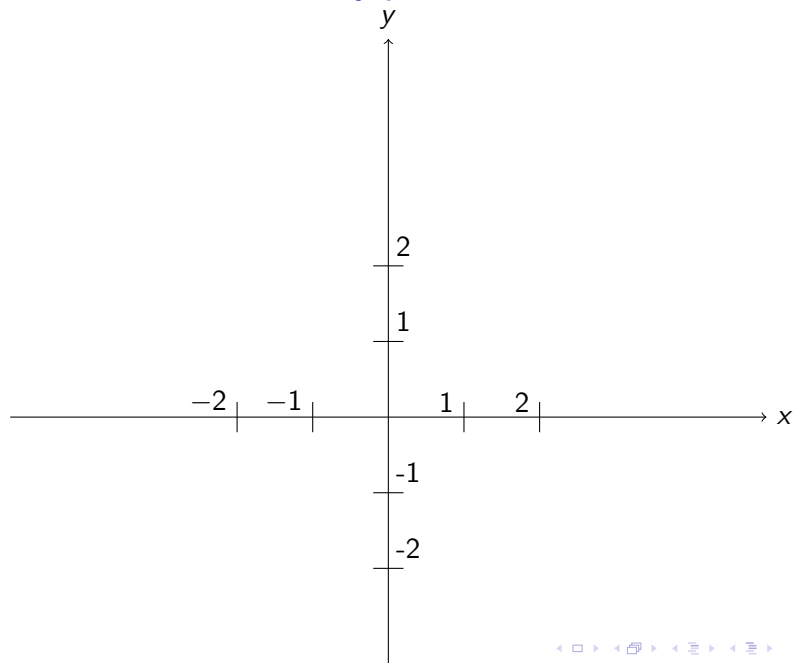
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Deletion-contraction rule:

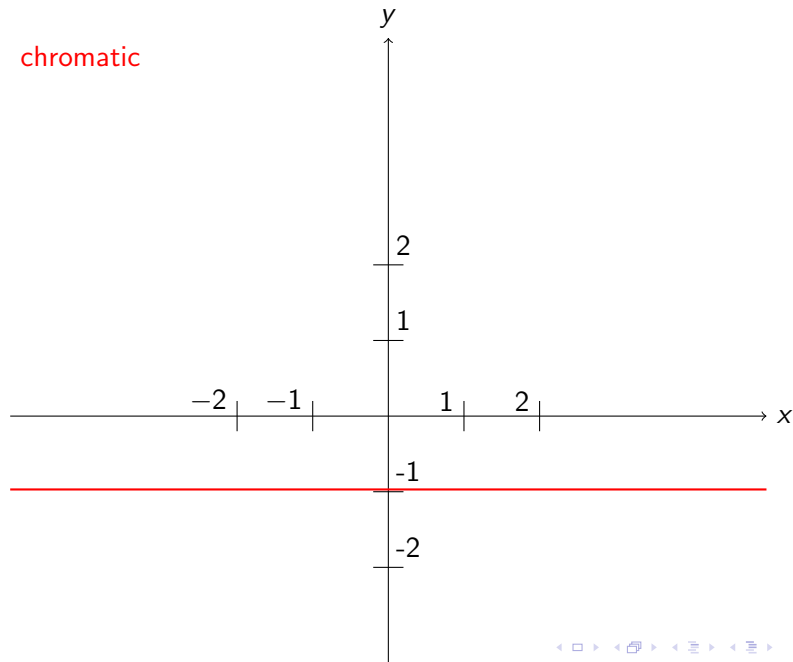
$$R(f; x, y) = x^{\log_2 \left(1 + \frac{f(\{e\})}{f(\emptyset)}\right)} R(f \setminus e; x, y) + y^{\log_2 \left(1 + \frac{\hat{f}(\{e\})}{\hat{f}(\emptyset)}\right)} R(f // e; x, y)$$

# Generalised Tutte-Whitney plane



# Generalised Tutte-Whitney plane

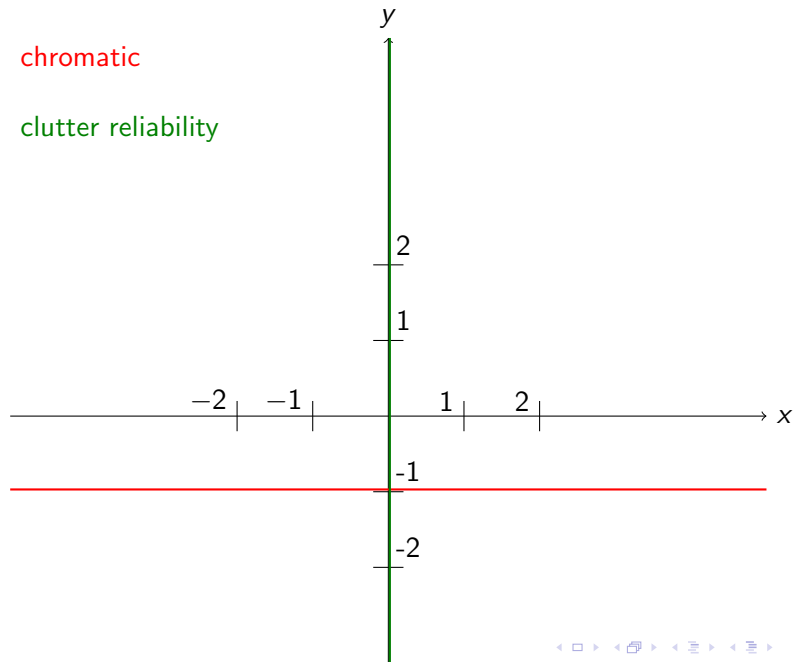
chromatic



# Generalised Tutte-Whitney plane

chromatic

clutter reliability

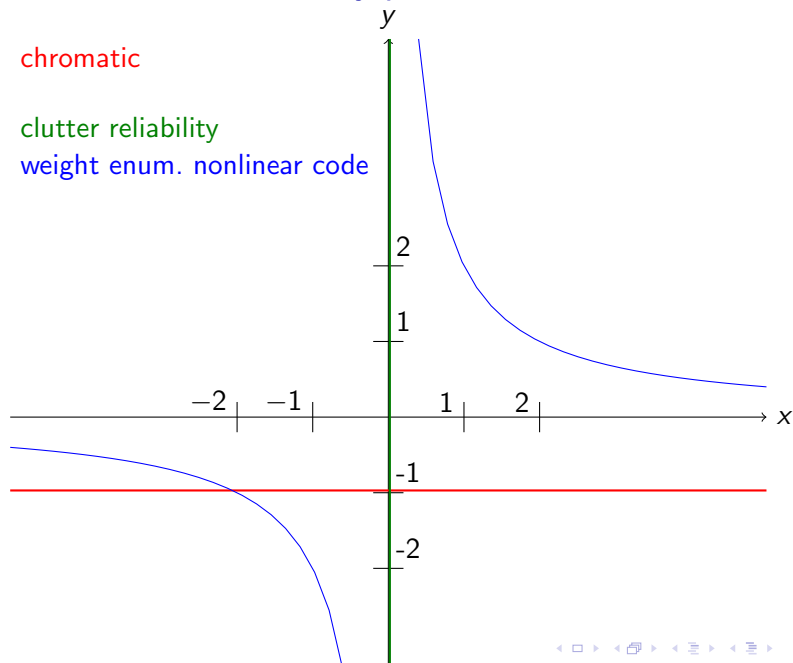


# Generalised Tutte-Whitney plane

chromatic

clutter reliability

weight enum. nonlinear code



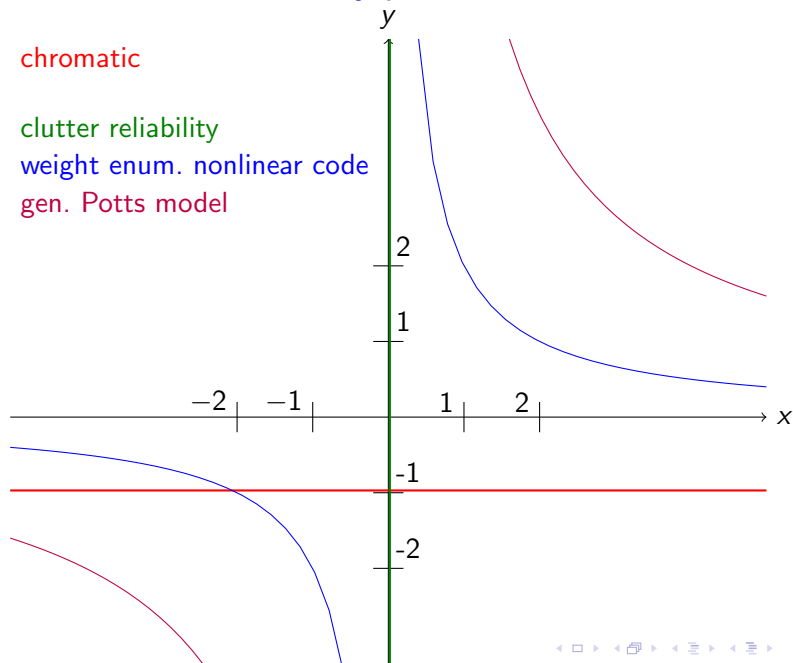
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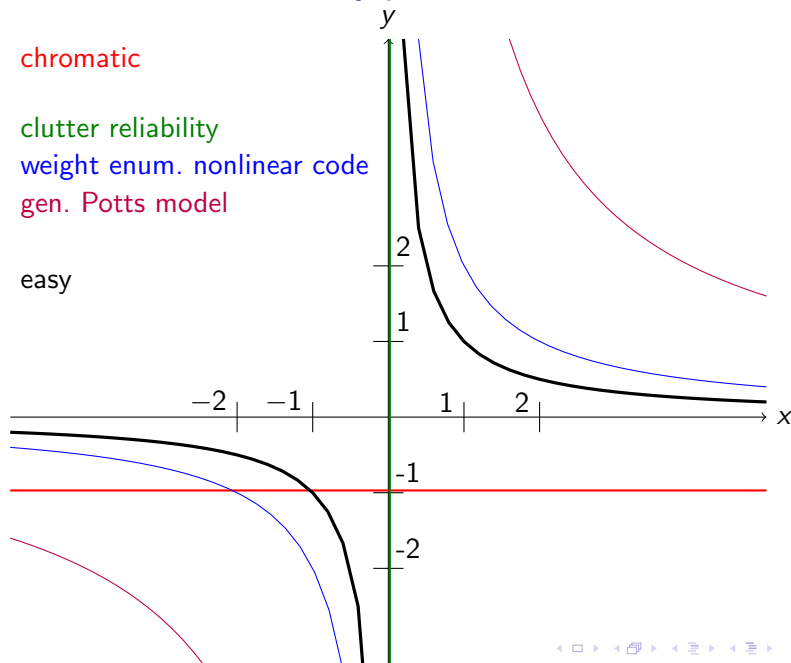
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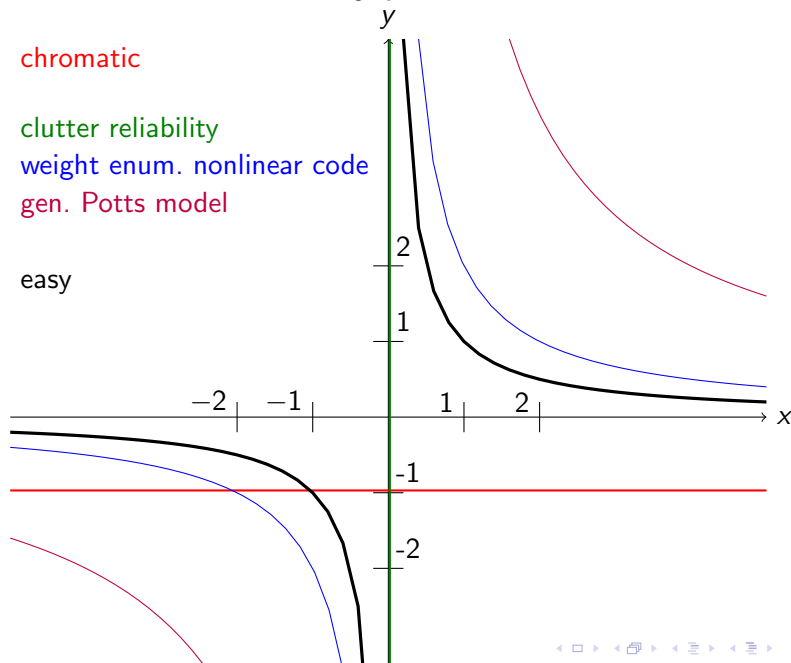
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# Generalised Tutte-Whitney plane

$y$

$x$

# Interpolating between contraction and deletion

For  $e \in E$ ,  $X \subseteq E \setminus \{e\}$ :

Contraction

$$(f // e)(X)$$

$$\frac{f(X)}{f(\emptyset)}$$

Deletion

$$(f \setminus e)(X)$$

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$\lambda$ -minor

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Fixed points:

$$\lambda = \pm\sqrt{2} - 1$$

## $\lambda$ -rank functions

Define  $Q^{(\lambda)}f$  by:

$$(Q^{(\lambda)}f)(W) = \log_2 \left( \frac{(1 + \lambda^*)^{|V|} \sum_{X \subseteq E} \lambda^{|X|} f(X)}{\sum_{X \subseteq E \setminus W} \lambda^{|W \cap \bar{V}|} (\lambda^*)^{|W \cap V|} f(X)} \right)$$

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$$(Q^{(\lambda)})^* = Q^{(\lambda^*)}$$

Inversion:

$$(Q^{\dagger(\lambda)}\rho)(V) =$$

$$(-1)^{|V|} (\lambda - \lambda^*)^{-|S|} \times$$

$$\sum_{W \subseteq V} (-1)^{|W|} (1 + \lambda^*)^{-|W|} (\lambda^*)^{|W \cap \bar{V}|} \lambda^{|W \cap V|} 2^{\rho(E) - \rho(E \setminus W)}$$

## A continuum of $\lambda$ -Whitney functions

$$R^{(\lambda)}(f; x, y) = \sum_{X \subseteq E} x^{Q^{(\lambda)}f(E) - Q^{(\lambda)}f(X)} y^{|X| - Q^{(\lambda)}f(X)}$$

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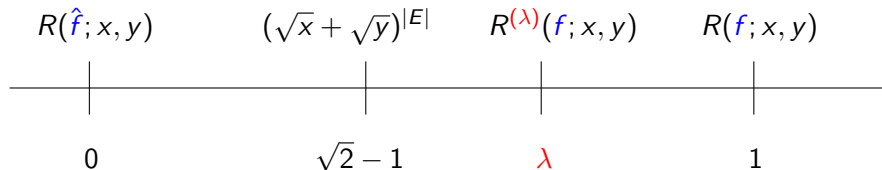
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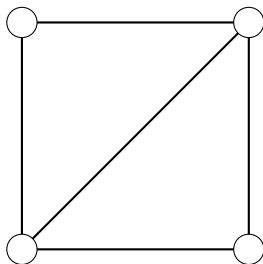
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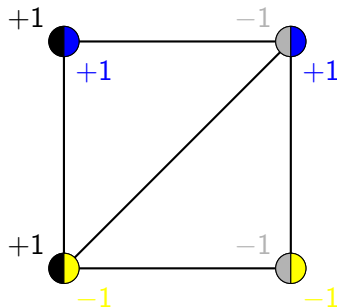
## Ashkin-Teller model (1943)

- 4-colourings (may be improper): colours are  $(\pm 1, \pm 1)$



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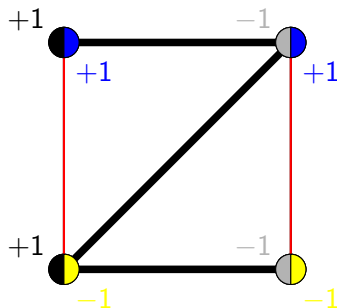
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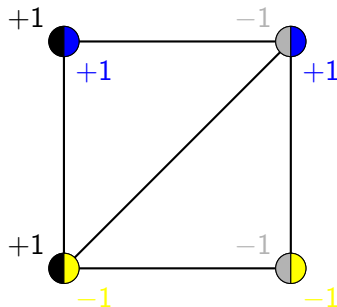
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*Left* colours:  
**Good** and **bad** edges

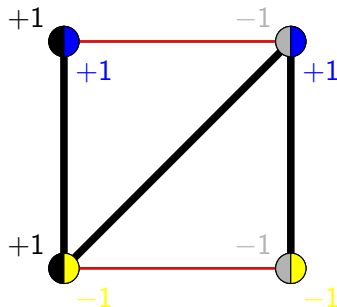
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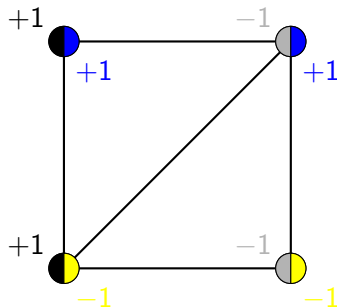
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*Right* colours:  
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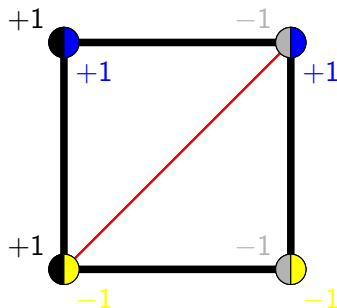
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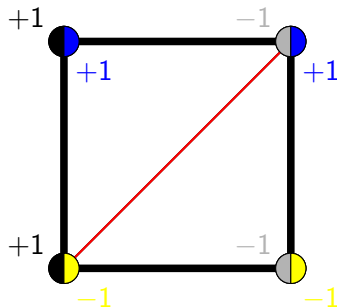
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**Good** and **bad** edges

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- 4-colourings (may be improper): colours are  $(\pm 1, \pm 1)$



*Product colours:*  
**Good** and **bad** edges

- Partition function (*symmetric* Ashkin-Teller):  
 $Z_{\text{AT}}(G; K, K', q) =$

$$e^{(2K+K')|E|} \sum e^{-\left( \begin{array}{l} K \cdot (\# \text{ good "left" edges}) \\ + K \cdot (\# \text{ good "right" edges}) \\ + K' \cdot (\# \text{ good "product" edges}) \end{array} \right)}$$

# Ashkin-Teller model (1943)

Special cases:

- ▶  $K = K'$ : Potts model (up to a factor)
- ▶  $K' = 0$ : product of two Ising models (each  $q = 2$ )

For these,  $Z_{\text{AT}}(G)$  is a specialisation of  $R(G : x, y)$ .

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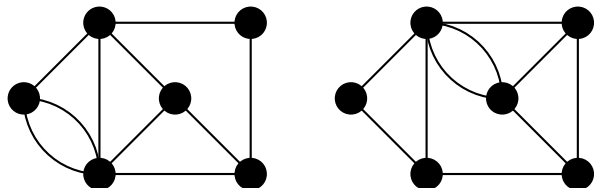
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In general,  $Z_{\text{AT}}(G)$  is *not* a specialisation of  $R(G : x, y)$ .

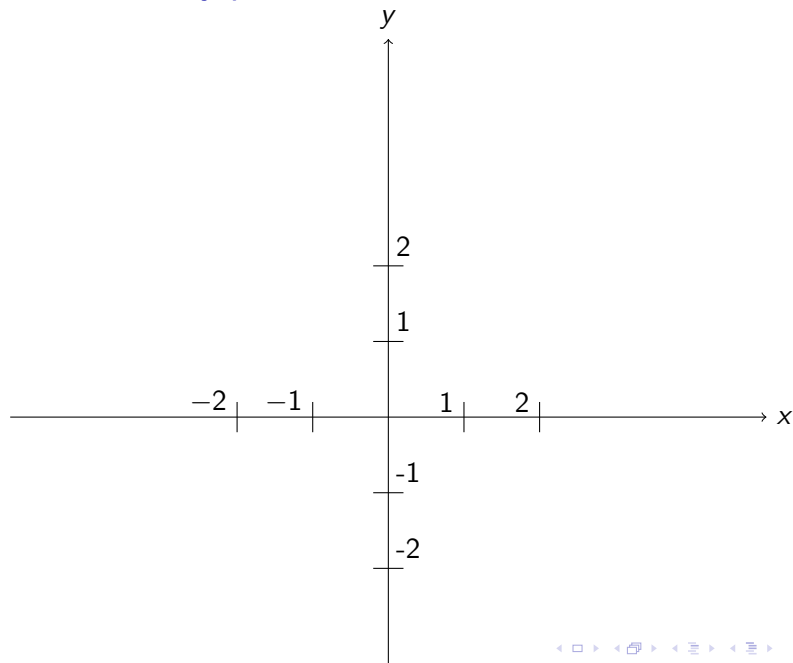
Example (M. C. Gray; see Tutte (1974)):



These graphs have same  $R(G; x, y)$ , but different  $Z_{\text{AT}}(G)$  (even in symmetric case).

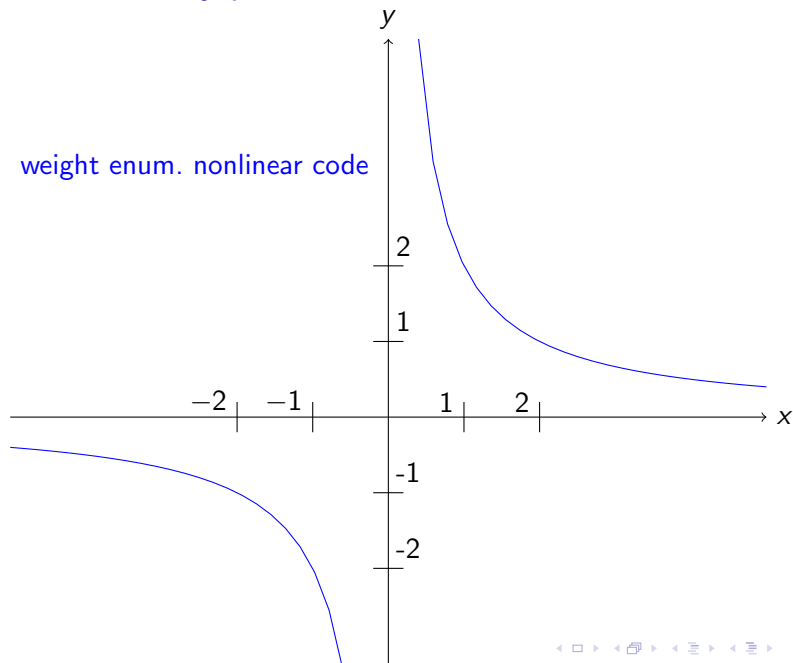


## $\lambda$ -Tutte-Whitney plane



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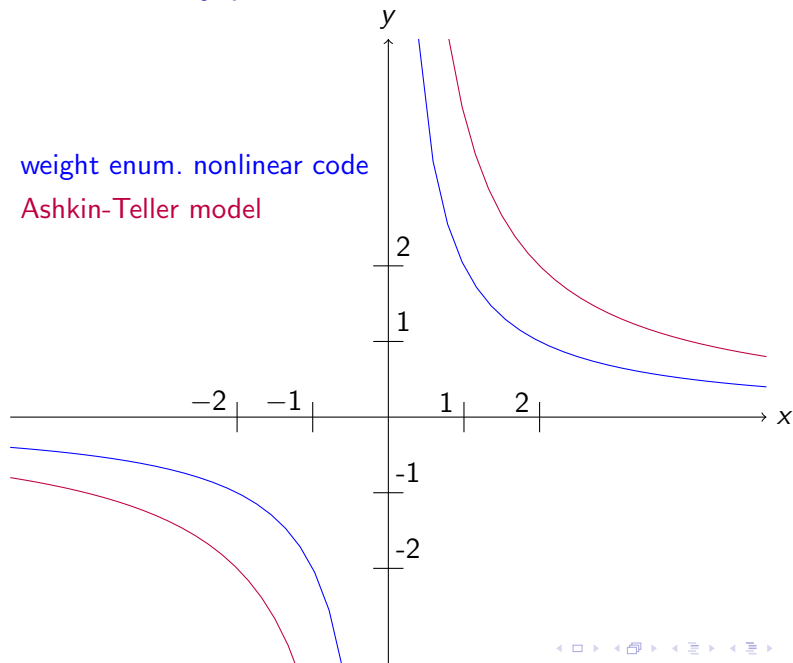
weight enum. nonlinear code



# $\lambda$ -Tutte-Whitney plane

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Ashkin-Teller model

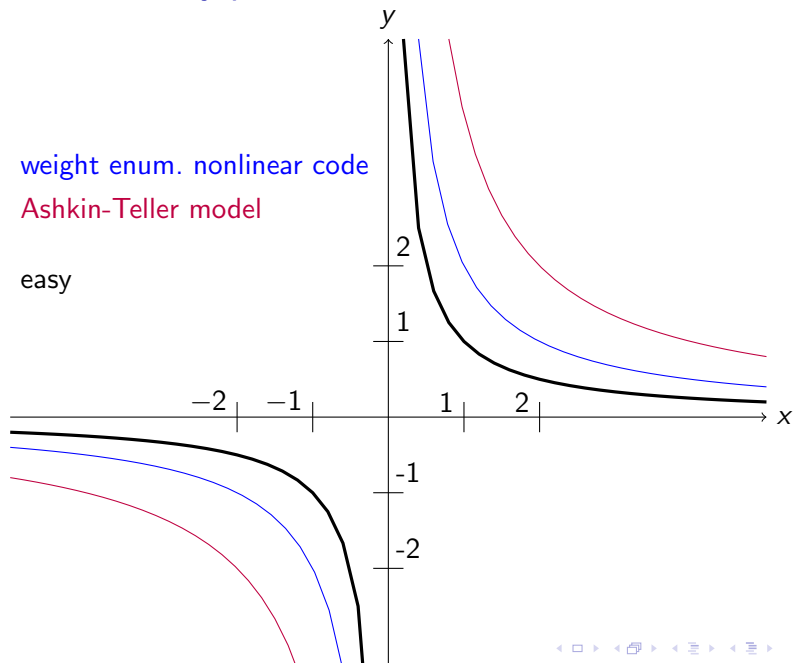


# $\lambda$ -Tutte-Whitney plane

weight enum. nonlinear code

Ashkin-Teller model

easy



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