An Introduction to Tries

Kevin Leckey

Monash University

21.09.2015

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 $\Xi_4 = 00000..., \quad \Xi_5 = 11111..., \quad \Xi_6 = 11100...$

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Task: Storage that allows fast search and insert/delete operations

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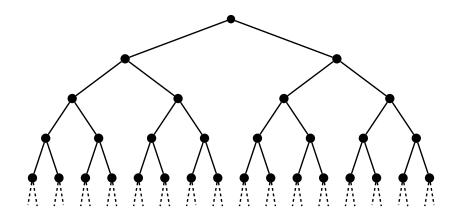
Task: Storage that allows fast search and insert/delete operations

→ Use tree-like data structures such as a Trie (Information retrieval)

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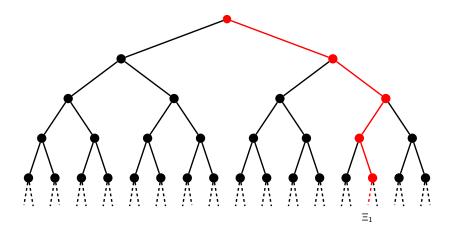
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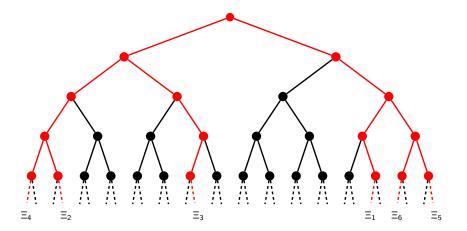
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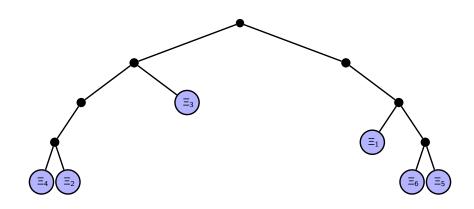
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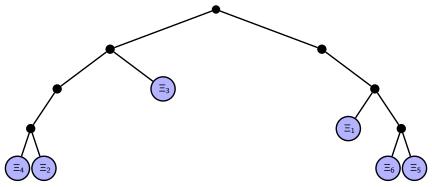


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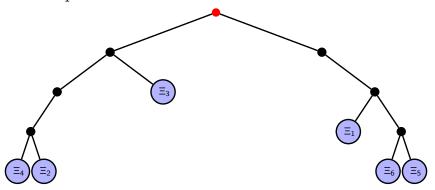
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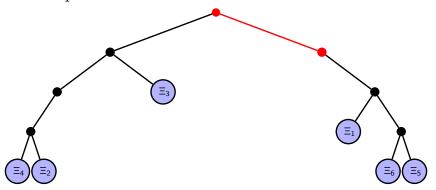
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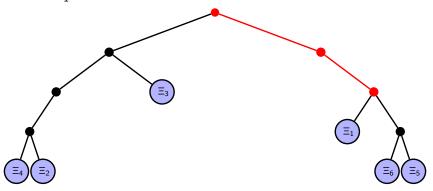
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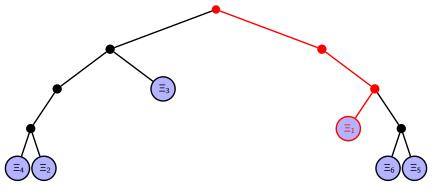
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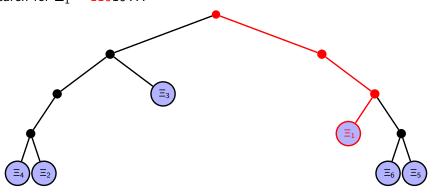
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Search for $\Xi_1 = {110}10\dots$

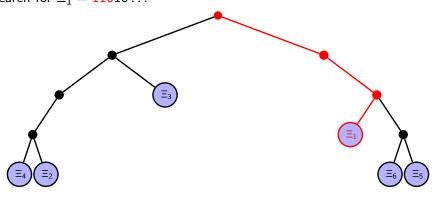


Search for $\Xi_1 = 11010...$



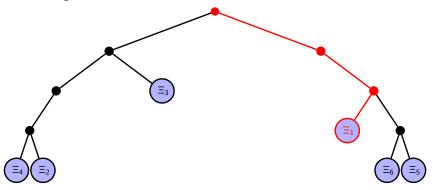
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Search for $\Xi_1 = {110}10\dots$



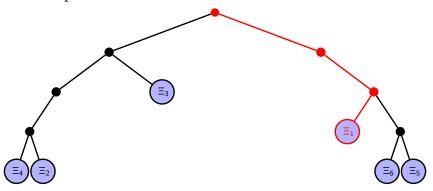
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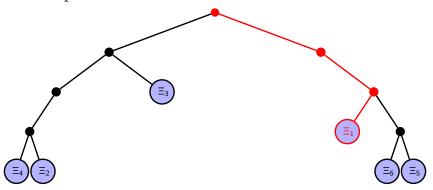
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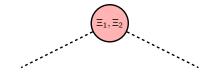
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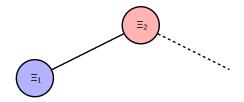
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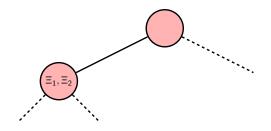
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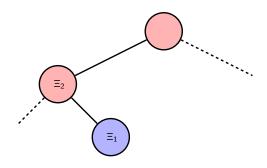
More general models allow ξ_1, ξ_2, \dots to be dependent (e.g. evolving as a Markov chain)

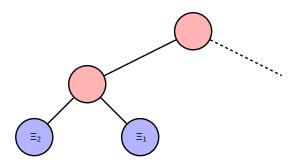


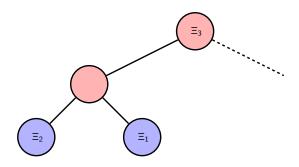


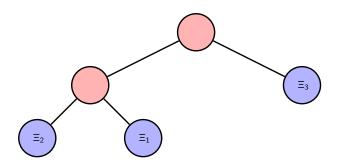


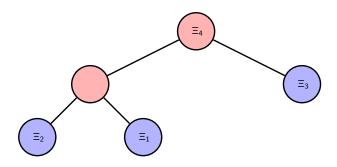


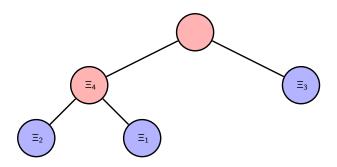


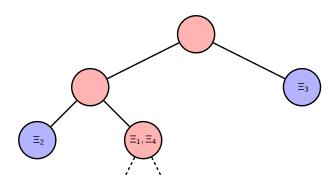


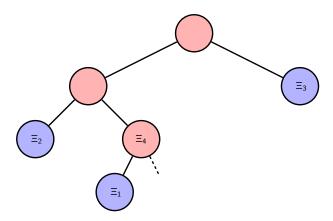


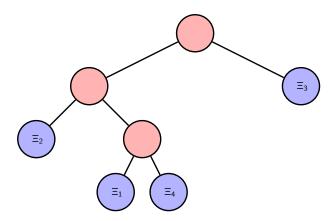


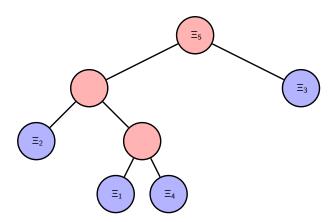


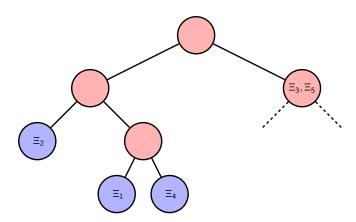


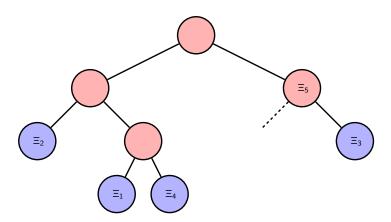


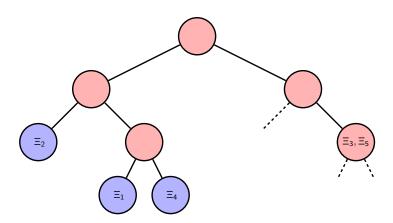


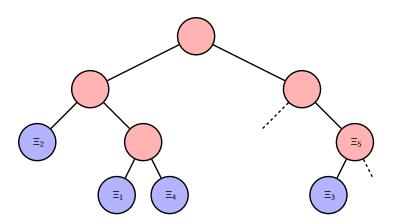


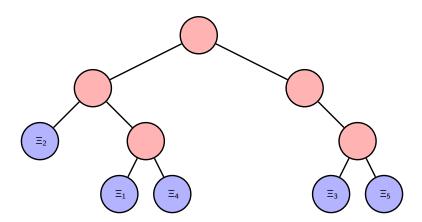












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$$\mathbb{P}(D_n \leq \alpha \log_2(n)) = (1 - n^{-\alpha})^{n-1} \stackrel{n \to \infty}{\longrightarrow} \begin{cases} 1, & \text{if } \alpha > 1, \\ 0, & \text{if } \alpha < 1. \end{cases}$$

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- Thm (Szpankowski '86): $Var(D_n) \sim \Phi(\log_2(n))$ with periodic function Φ

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Consequence: $\mathbb{P}(H_n > \alpha \log_2(n)) \to 0$ for $\alpha > 2$

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Thm (Regnier '82):

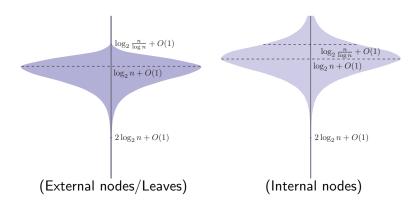
$$\mathbb{E}[H_n] \sim 2\log_2(n) \qquad (n \to \infty)$$

(Flajolet, Steyaert '82 \rightarrow periodic second order term)

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Profile (Park, Hwang, Nicodème, Szpankowski):

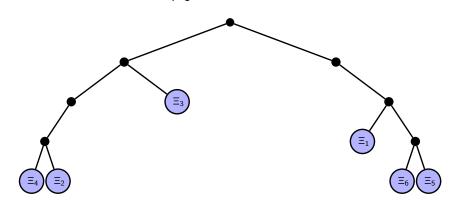


Consider *n* words Ξ_1, \ldots, Ξ_n . External Path Length:

$$L_n:=\sum_{i=1}^n D_{n,i}, \qquad D_{n,i}=D_n(\Xi_i).$$

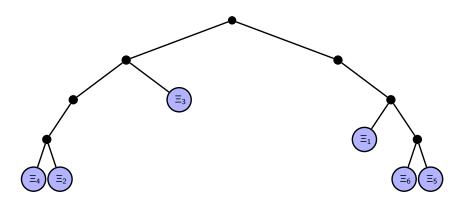
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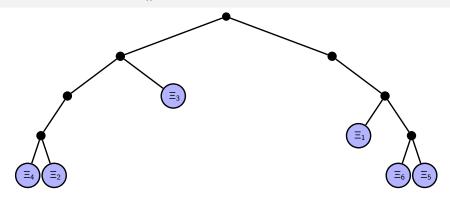
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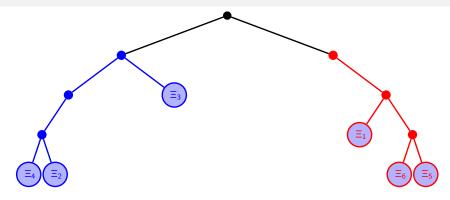
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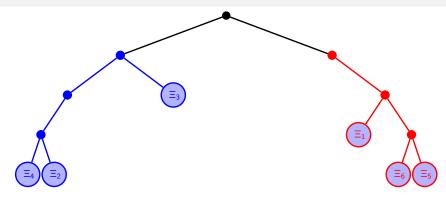


Example: $L_6 = 2 + 3 + 4 \cdot 4 = 21$

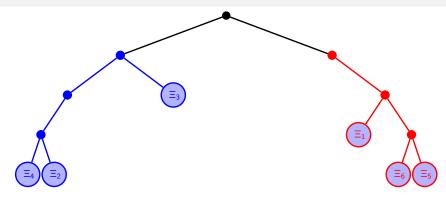
A Recursion for L_n





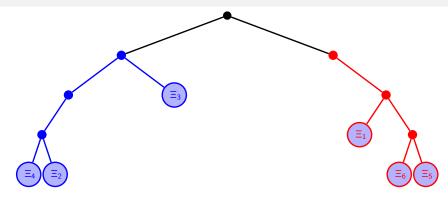


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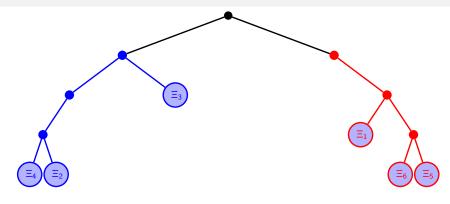
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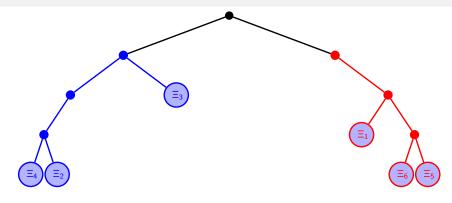
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2. Find the Limits: $(A_{n,1}, A_{n,2}, b_n) \longrightarrow ???$

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- Input model not very realistic, what about more general input models?

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Even more general: Dynamical Sources Model by Vallée

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$$\mathcal{H} = \pi_0 \left(-p_{00} \log(p_{00}) - p_{01} \log(p_{01}) \right) + \pi_1 \left(-p_{10} \log(p_{10}) - p_{11} \log(p_{11}) \right)$$

with stationary distribution

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Depth for Markov Sources:

$$\mathbb{E}[D_n] \sim \frac{1}{\mathcal{H}} \log(n)$$

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- Lempel-Ziv Parsing Scheme (data compression)