How To Determine If A Random Graph With A Fixed Degree Sequence Has A Giant Component

Bruce Reed

Monash School of Mathematical Sciences Colloquium December 3rd, 2015

Looking for Clusters I: **Epidemiological Networks**

Epidemic spread in networks: Existing methods and current challenges

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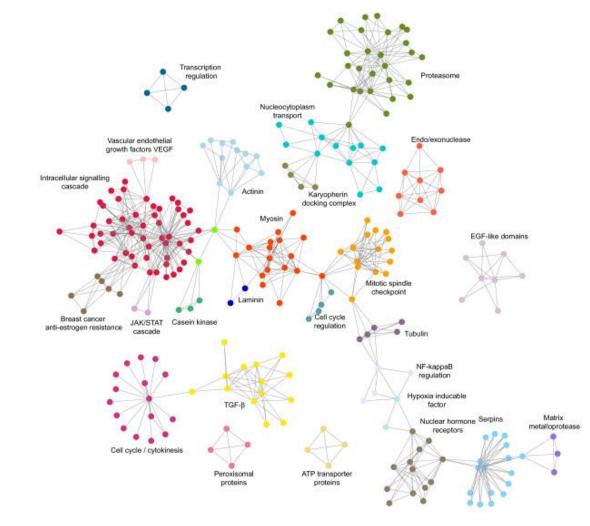
March 11, 2014

Abstract

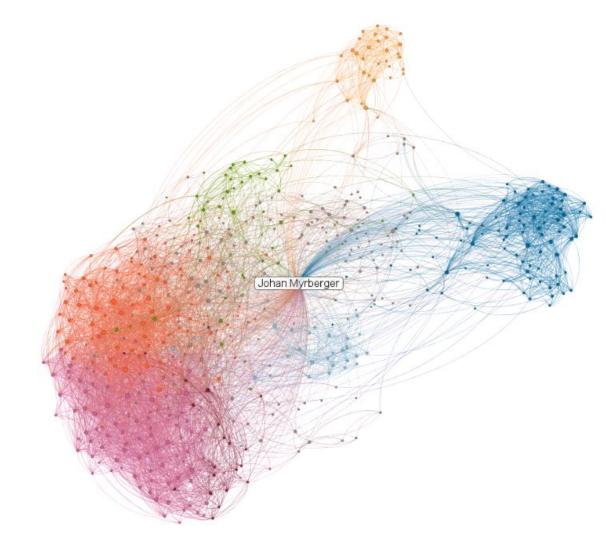
We assume that the disease spreads through contacts and infected individuals recover into an immune state. We discuss a number of existing mathematical models used to investigate this system, and show relations between the underlying assumptions of the models. In the process we offer simplifications of some of the existing models. The distinctions between the underlying assumptions are subtle, and in many if not most cases this subtlety is irrelevant. Indeed, under appropriate conditions the models are equivalent. We compare the benefits and disadvantages of the different models, and discuss their application to other populations (e.g., clustered networks). Finally we discuss ongoing challenges for network-based epidemic modeling.

Looking for Clusters II: **Communication Networks**

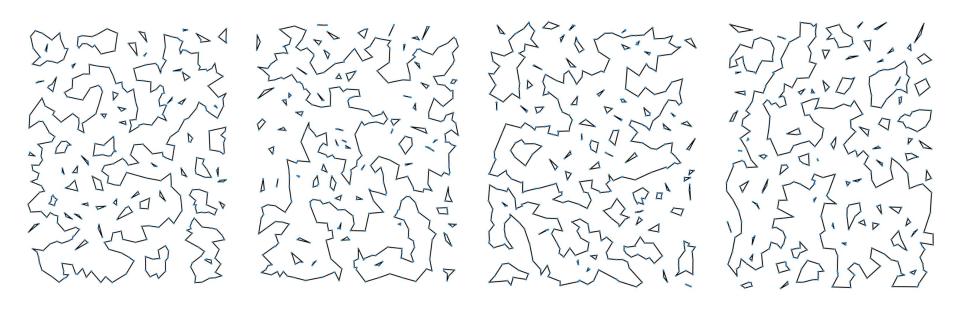
Looking for Clusters III: Biological Networks



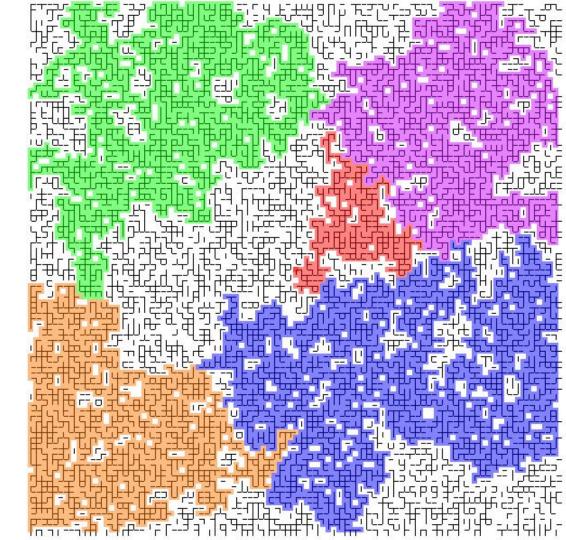
Looking for Clusters IV: Social Networks

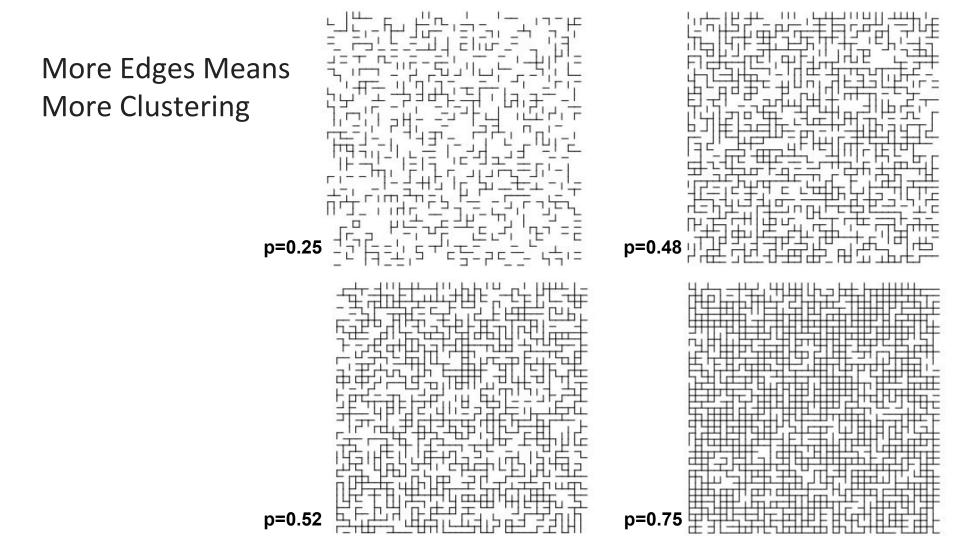


Looking for Clusters V: Euclidean 2-factors

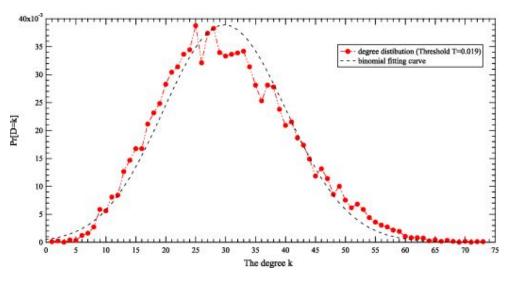


Looking for Clusters VI: Percolation

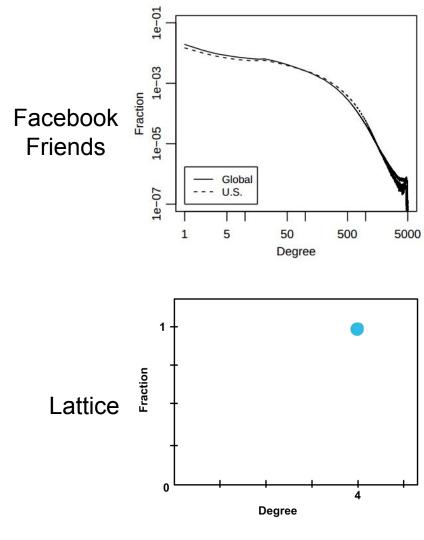




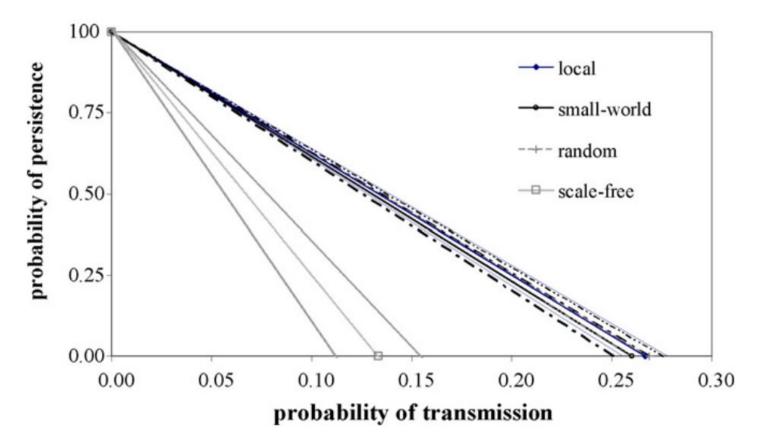
Degree Distributions Differ



Classic Erdős-Renyi Model



Network Structure Affects Cluster Size



Random Networks as Controls

A common technique to analyze the properties of a single network is to use statistical randomization methods to create a reference network which is used for comparison purposes.

Mondragon and Zhou, 2012.

Does a uniformly chosen graph on a given degree sequence

have a giant component?

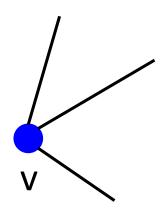
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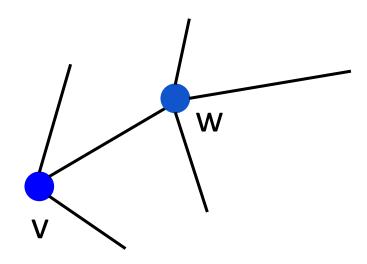
For a sequence D of nonzero degrees, G(D) is a uniformly chosen graph with degree sequence D.

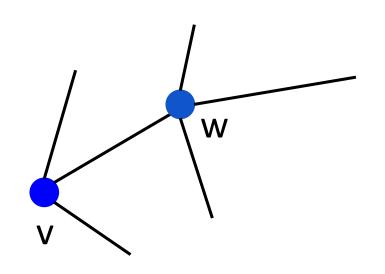
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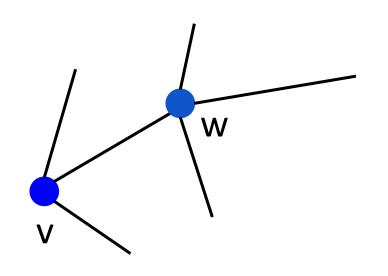
Will assume D is non-decreasing and all degrees are positive.





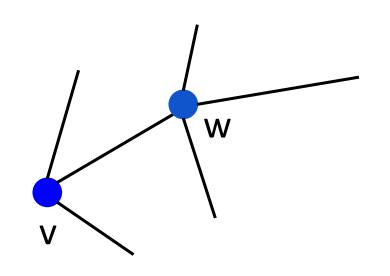


Change in number of open edges: d(w) - 2



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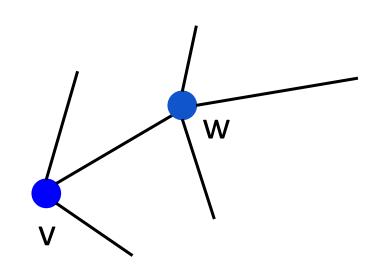
Probability pick w: $d(w) / \sum_{u} d(u)$



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Giant Component if and only if $\sum_{u} d(u)(d(u)-2)$ is positive??

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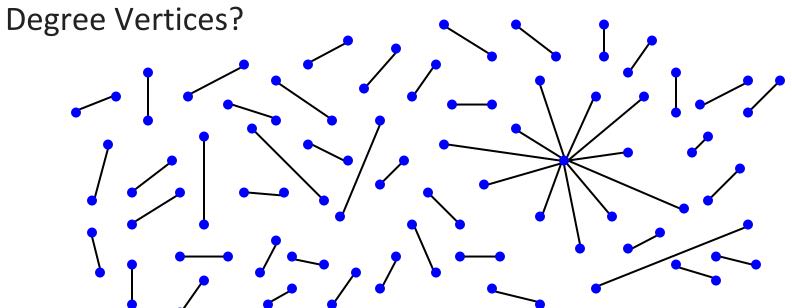
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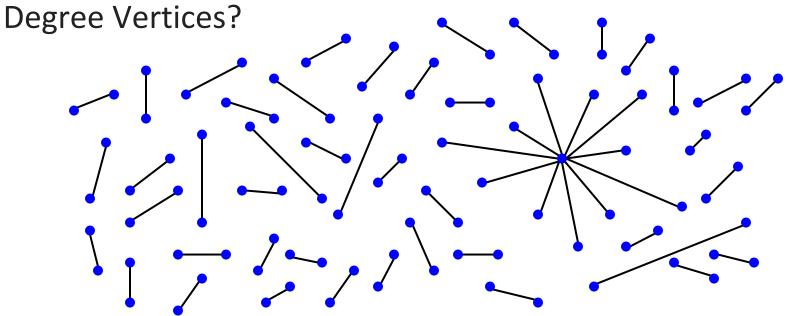
Molloy-Reed(1995) Result

Under considerable technical conditions including maximum degree at most $n^{1/8}$:

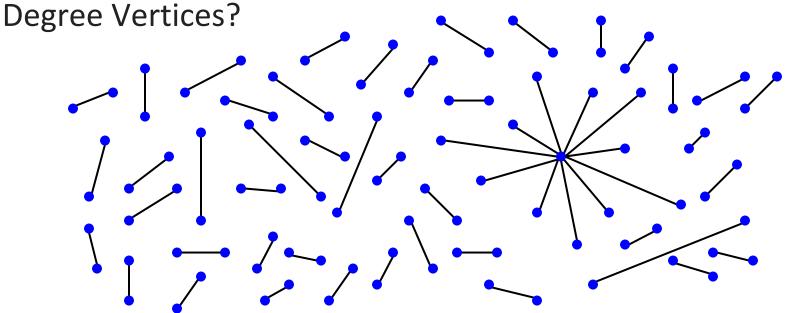
$$\sum_{u} d(u)(d(u) - 2) > \varepsilon n \qquad \text{implies a giant component exists.}$$

$$\sum_{u \in \mathcal{U}} d(u)(d(u) - 2) < -\varepsilon n \quad \text{implies no giant component exists.}$$

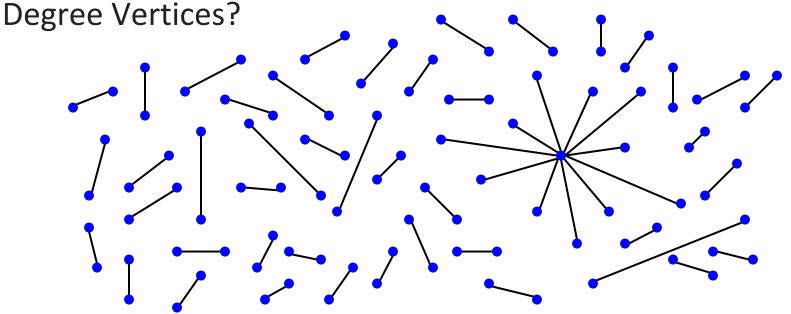




Because it is false.



Cannot translate results from the non-simple case.



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Hard to prove concentration results.

OUR QUESTION REVISITED

Does a uniformly chosen graph on a given degree sequence have a giant component?

For a sequence D of nonzero degrees, G(D) is a uniformly chosen graph with degree sequence D.

Will assume D is non-decreasing and all degrees are positive.

Four Definitions

M is the sum of the degrees in D which are not 2.

D is f-well behaved if M is at least f(n).

$$j_D = \min (i \text{ s.t. } \sum_{j=1}^i d_j(d_j - 2) > 0, n)$$

$$R_D = \sum_{j=0}^{n} d_j$$

One Crucial Observation

$$\sum_{j=1}^{n} d(u)(d(u)-2) \text{ is at least R}_{D}$$

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and for some V > 0 remains above $R_D/2$ until the sum of the degrees of the vertices explored is at least VR_D .

But goes negative once all the vertices with index $> j_D$ are explored.

Two Theorems

Theorem 1: For any $f \to \infty$ and $b \to 0$, if a well behaved degree distribution D satisfies $R_D \le b(n)M$ then G(D) has no giant component

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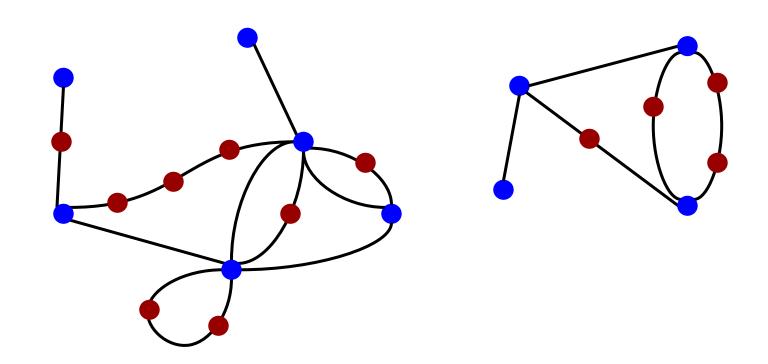
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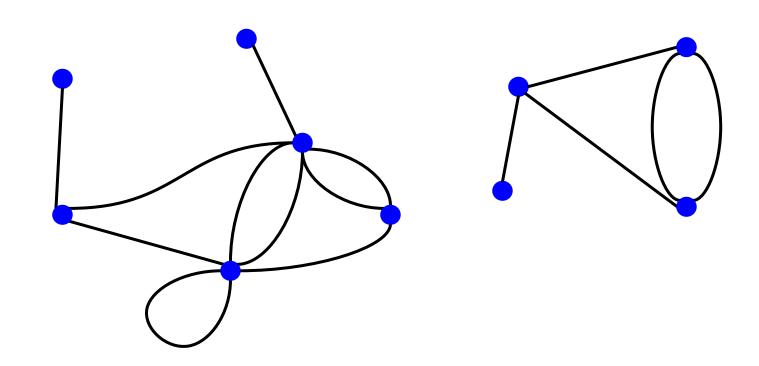
Theorem 2: For any $f\to\infty$ and $\epsilon>0$ if a well behaved degree distribution D satisfies $R_D \geq \epsilon M$ then G(D) has a giant component

(Joos, Perarnau-Llobet, Rautenbach, Reed 2015)

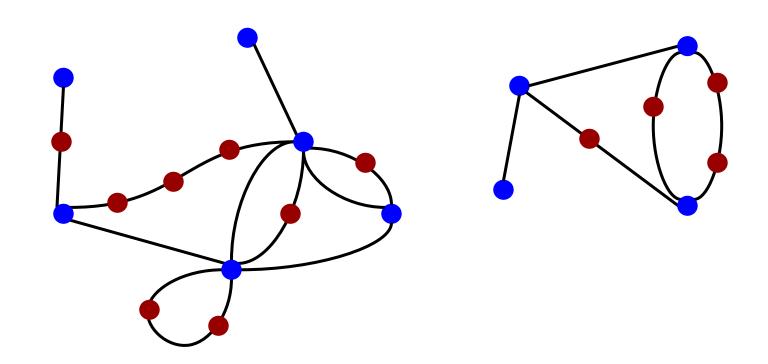
Why we focus on M and not n

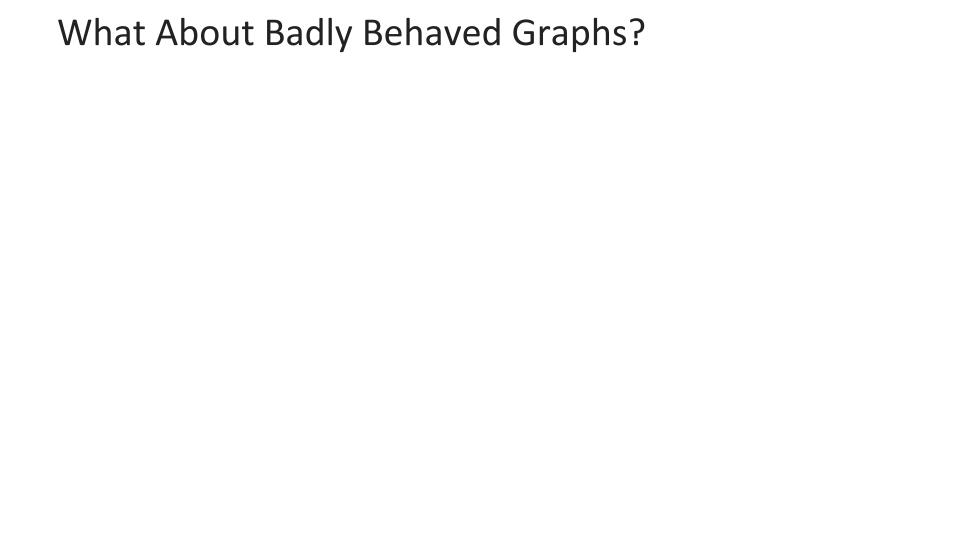


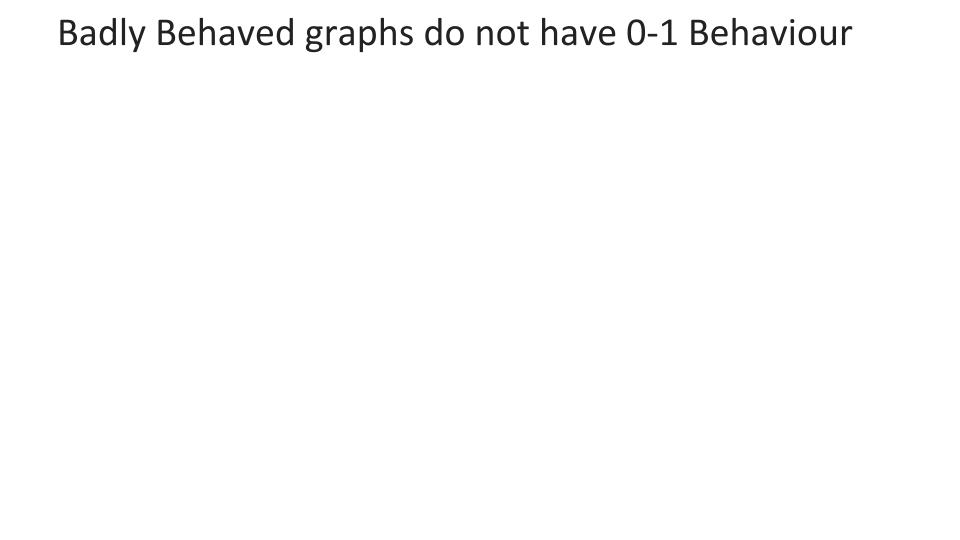
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Badly Behaved graphs do not have 0-1 Behaviour

For all 0<ε<1, the probability of a component of size at least εn lies between c and 1-c for some constant c between 0 and 1.

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If all vertices of degree 2 just taking a random 2-factor.

If M is at most some constant b, with probability p(b)>0 all but $\epsilon n/2$ of the vertices lie in cyclic components.

Two Theorems

Theorem 1: For any $f \to \infty$ and $b \to 0$, if a well behaved degree distribution D satisfies $R_D \le b(n)M$ then G(D) has no giant component.

Theorem 2: For any $f\to\infty$ and $\epsilon>0$ if a well behaved degree distribution D satisfies $R_D \geq \epsilon M$ then G(D) has a giant component

(Joos, Perarnau-Llobet, Rautenbach, Reed 2015)

Differences in the Proof

Determine if there is a component K of the multigraph obtained by suppressing degree 2 vertices satisfying:

(*)
$$|E(K)| > \varepsilon'M$$
.

Use a combinatorial switching argument to obtain bounds on edge probabilities in this multigraph.

Differences in the Proof - When No Giant Component Exists

Begin the random process with a large enough set of high degree vertices that our process has negative drift.

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Show drift becomes more and more negative over time, if the process does not die out.

Differences in the Proof - When A Giant Component Exists

Focus on the set $H = \{v \mid d(v) > (\sqrt{M})/log(M)\}$

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We can show, using our combinatorial switching argument, that depending on the sum of the sizes of the components intersecting H, either

- (a) there is a giant component containing all of H, or
- (b) we can reduce to a problem with H empty.

Demonstrating The Switching Argument

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Theorem: If |E|>8n log n then,

Prob(G has a component with (1-o(1))n vertices)= 1-o(1).

Future Work

Tight bounds on the size of the largest component in terms of $R_{\rm D}$

Thank you for your attention!

